Error Correction (Shor) - page 1

x107+ B122

Thursday, November 29, 2018 9:46 PM Classical error correction is accomplished try encoding data in a redundant way so that if part of the date is lost, the problem can be detected and then corrected. Challenging in the guantum setting because: 1) Data can how be copied (No Cloning) 2) Meconement destroys grantine information. 3) Erroro can be continuous. For a single qubit, one possible emoris a bit flip:  $|0> \rightarrow |1>$ ,  $|1> \rightarrow |0>$ . More generally: Xlo>+ fl=> > flo>+ xl=> We will use a standard repitition code: 14> 14>= xlo>+ pl2> 10> 10 (14>100) -> 10> 10> (14>100) -> xlooo>+pl111>  $\begin{array}{c} |0\rangle \rightarrow |000\rangle \\ |1\rangle \rightarrow |11\rangle \end{array}$ ( xloo>+ pl11>) @ lo> ( 2 0007 1 B/117)

Error Correction (Shor) - page 2 2 hx>/x| T= Thursday, November 29, 2018 9:46 PM -P We can use the following projectors to detect an error: P1 1000> = 0 000 > 1000 + 111 > 111  $P_0 =$  $P_1 = P_2 =$ P2 010> = 1010> 100><100 + 1011>Lo11 + 1010><010] + 101><101] P3 ( d 1001>+ f 1110>) 1001720011 + 1110721101  $P_{\mathbf{z}} =$ We can store the result in an anxiliary register: 1-> 1000> data Syndrome 1000> + B 111> Can do  $P_1 \otimes XII + (I-P_1) \otimes I^3$   $(\alpha | 100 > + \beta | 011 >) | 000 > - (\alpha | 100 > + \beta | 011 >) | 100 >$ Note that measuring the syndrome does not destroy the gnawnum information in the first register. It just indicates that an error has occurred. Then we can do a conholled X to correct the error : 1st bit of syndrome is control 1st bit of date is the target This works even if error happens in Superposition. Flip bit 1 w/ amplitude /15 to (XII + IXI) Plip bit 2 w/ amplitude /152

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定 (x/100>+ B/011>)/100>+ 売 (x/010>+ B/101>)/010> After even correction: (x 1000>+ plin>) tz (100>+ 1010>) Can measure syndrome before or after. This method does not work with 2 or more errors => Only O or I errors. If each gubit is flipped independently with probability p, the probability that this method workers is:  $(1-p)^{3} + 3p(1-p)^{2} = 1 - 3p^{2} + 2p^{3}$ No mors 1 erm. (probernon in 15t qubot may is p(1-p)2). This method is an improvement over he ever correction M:  $3p^2 - 2p^3 > p = p < \frac{1}{2}$ 

Friday, November 30, 2018 8:50 AM An alternative view : Instead of performing measureness P. P. P. B Do two measurements: ZZZI and IZZZ3 ZZI = ( 100×00 + 111><11)@I - (101×01 + 10><101)@T This has eigenvalues ± 1 and gives 1 bit of information. ZIZZI and IZZZ together give 2 bits of information. Zitz Compares quibits 1 and 2. Cigenvalue is +1 if gubits 1+2 are equal. Cigenvalue is -1 if gubits 1+2 are defined Note The two bits of information fell you if I where a singer bit flip occurred! 0 |-|-| 2125 J227 ho flip. + \ +1 No information about d'or p is obtained. 312 + ( - | 18-+ \ -1 24 -1 -1 The syndrome is the two tits resulting from this measurement.

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Another possible ever is a phase flip: with probability p, Z operator is applied to a gubit. xlo>+Bl2> -> xlo> - Bl1>. Can turn a phase flip channel into a bit flip channel Z flips (+> and 1-> Encode 10> as 1+++> Encode 11> as 1---> Analysis is the same as with the bit flip Chancel. Measurements done in the Ity I-> basis. Syndrome measurement 10> - $H^{\otimes 3} = X_1 X_2$   $H^{\otimes 3} = X_2 X_3$ Recovery on the first gubit is done by H, X, H, = Z1 For example, if syndrome indicated har 1+++> becaue 1-++> then Z1 1-++> = 1+++>

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The Shor Code Uses 9 qubits to encode a Single gubit and protect agained any error as long as it effects only a single gubit of the code. First encode the qubit using the phase-flip code then encode each of the gubits using the bit flip code. This hellod of encoding using hestel levels is Called concatenation and is useful for combining benefits of different codes.

Friday, November 30, 2018 8:50 AM (1000>+111)(1000>+111)(1000>+111) 0> 1000/-1111) (1000/-1111) (1000/-111) 12) 252 14> H 107 10> 107 10> 10) 107 10> To connet for tit flips, use the error correction procedure for tit flips in each block separately: ZIZZ + ZZZZ -> connoch ZyZs + ZsZb - connel ZyZs + ZsZg - connel

Friday, November 30, 2018 10:08 AM (1000>+1111>)(1000>+1111>)(1000>+1111>) 252. 0 > ( 1000>-1111>) (1000> - 1111>) (1000>-1111) 1) Now correct for phase flips between blocks Measure: X1X2X3X4X5X6 Xy X5X6 X1 X8X9 For example suppose there is a phase flip in the first block: (1000>-111)(1000>+1111>)(1000>+1111>) 0 > 25 ( 1000>+ 1111>) (1000> - 1111>) (1000> - 1111>) 1) Measuring X1X2X3X4X5X6 will yield -1 To corner, apply ZIZZZ3 d ( 1000> - 1111>) + f ( 1000>+ 1111>) 

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If the error is a bat flip and phase flip on the same gubit X1Z1 lach correction procedure (bit flip then phase flip) will work independently. Any 1-gubit ever can be appressed as: E= lo I + l, X, + l2 Z, + l3 XZ, E 4> = eo 14> + e, X1 14> + e2 Z1 14> + e3 X1Z1 14> Meesning the syndrome collapses to one of the four states: 14> X,14> Z,14> X,Z,14> and the appropriate to covery can be applied to get back to 147. A whole continuum of errors can be handled by only correcting a discrete set.