

# What have we seen?

- ▶  $\mathbb{N} = \{0, 1, 2, \dots\}$
- ▶  $|\mathbb{N}| = \aleph_0$
- ▶ There are this many odds, integers ( $\mathbb{Z}$ )
- ▶ There are this many triples of naturals
- ▶ There are this many rational numbers

# And now $\mathbb{R}$

- $\delta(k, x)$  :  $k$ th digit of  $x$  after decimal point

*FSOC* ► ex:  $\delta(2, \pi) = 4$  *# I'm thinking of*

- I claim  $f : \mathbb{N} \rightarrow \mathbb{R}$  is bijective.
- To disprove: find  $r \in \mathbb{R}$  s.t.  $\exists y f(y) = r$   
 is  $\lfloor f(0) \rfloor = 6$  ? If so,  $\lfloor r \rfloor = 7$   
 else  $\lfloor r \rfloor = 6$

for all  $k \in \mathbb{N}, k \geq 1$ ?

set  $\delta(k, r) = 7$  if  $\delta(k, y) = 6$   
 else set = 6

# Cantor's Theorem

**Claim:**  $|\mathcal{P}(A)| > |A|$

Suppose FSOC it is not. Then there is some  $f : A \rightarrow 2^A$  that is surjective.

$$\text{def } B = \{x \in A \mid x \notin f(x)\}$$

$$B \subseteq A \text{ so } B \in \mathcal{P}(A)$$

$$B \in 2^A$$

$$\exists b \quad f(b) = B ?$$

is  $b \in B$ ?

Yes?  $\rightarrow \leftarrow$

No?  $\rightarrow \leftarrow$

CompSci 162  
Spring 2023 Lecture 3:  
Formal Languages, Automata

# Five Languages

$L_1 = \{a, abb, aaaa\}$  finite

$L_2 = \{a^n \mid n \in \mathbb{N} \text{ is prime}\}$

$L_3 = \{b^n a^n b^m \mid n, m \in \mathbb{N} \text{ and } n \equiv m \pmod{3}\}$

$L_4 = \text{The set of all } w \in \Sigma^* \text{ with at most three a's}$

$L_5 = \{a^n \mid n \in \mathbb{N} \text{ and } \exists x, y, z \in \mathbb{N} - \{0\} \text{ such that } x^n + y^n = z^n\} = \{a, aa\}$

all possible strings over an alphabet

**Question 9** Which of these languages are infinite?

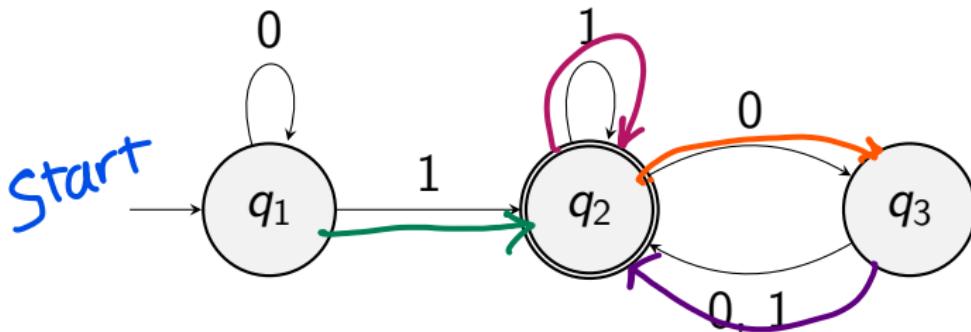
## Question 10

Are all languages (over a given alphabet) defined by some string in a suitable meta-language?

See Ed  
discussion

# Question 11

What happens when input is “1101”?



$\bigcirc$  = "accept" state

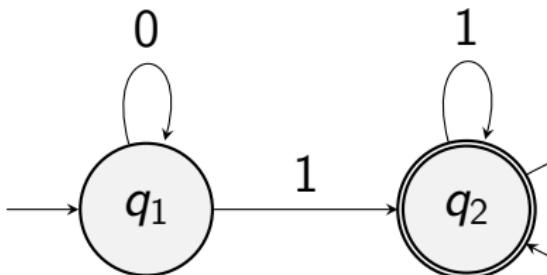
$\bigcirc$  = "reject" state (implicit)

This machine accepts 1101

# Formal Definition

$Q$  = set of states

$q_0$  = start state



$$Q = \{q_1, q_2, q_3\}$$

$$q_0 = q_1$$

$$\delta : Q \times \Sigma \rightarrow Q :$$

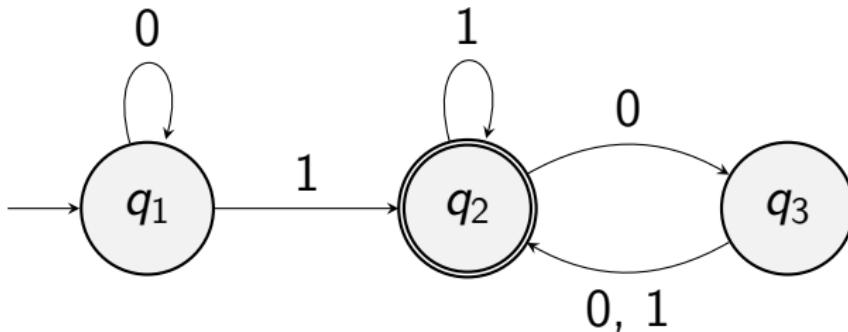
$q_1$	$q_1$	$q_2$
$q_2$	$q_3$	$q_2$
$q_3$	$q_2$	$q_2$

input alphabet

$$\Sigma = \{0, 1\}$$

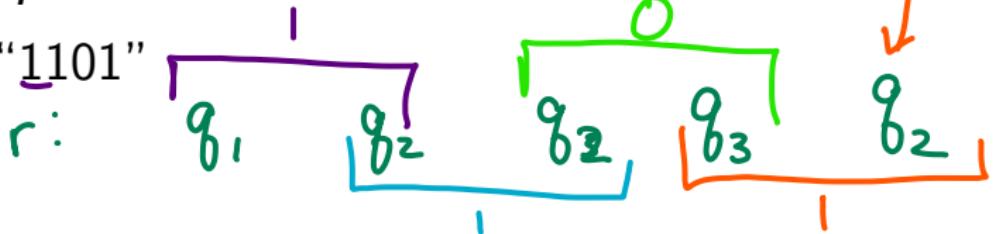
$$F = \{q_2\}$$

# Formal Definition of Computation



- $r_0 = q_0$
- $\delta(r_i, w_{i+1}) = r_{i+1}$  for  $i = 0 \dots n - 1$
- $r_n \in F$

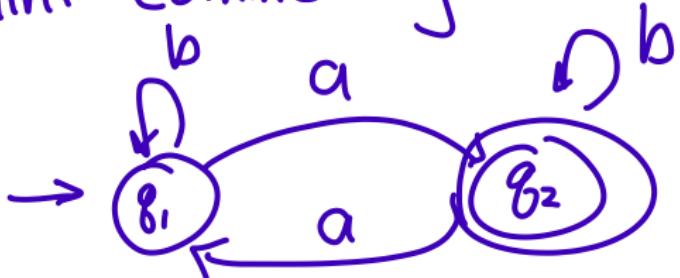
Input = "1101"



## Question 13

Design a DFA over the language  $\Sigma = \{a, b\}$  that accepts all strings with an odd number of instances of the letter  $a$ .

hint: comment your "code"!



//  $q_1$ : even #as

$q_2$ : odd # as