

1. In lecture, you saw the Path Property of MSTs, but we did not prove it. We will do so here. For any two vertices  $x$  and  $y$ , the path from  $x$  to  $y$  through the minimum spanning tree has the minimum possible weight for its heaviest edge. This is the **path property** of Minimum Spanning Trees.
2. Suppose we have a weighted graph  $G$  (undirected, connected,  $n$  vertices,  $m$  edges) with a unique minimum spanning tree (you may assume distinct edge costs too). Our goal is to compute the *second-best* spanning tree – that is, a spanning tree of  $G$  for which the only other spanning tree with a lower cost is the minimum spanning tree.
  - (a) Let's say that two spanning trees differ by  $k$  edges if each one includes exactly  $k$  edges that the other doesn't. In other words, they have exactly  $(n - 1) - k$  edges in common. There must exist a second-best spanning tree that differs from the MST by only 1 edge. Try to explain (or even prove!) why this fact is true.
  - (b) With that fact in mind, design an algorithm to find a second-best spanning tree in a graph.