

CompSci 161
Winter 2023 Lecture 11:
Dynamic Programming II:
Longest Common Subsequence

Longest Common Subsequence

Input: Two sequences (strings etc)

Output: Longest common subsequence.

Examples of common subsequences:

exercise

determine

eerie

morning

triangle

ring

toward

thousand

toad

LCS: Recursive Solution

Cases: $O(mn)$
time each: $O(1)$

Let $LCS(n, m)$ be the length of the longest common subsequence of $X[1 \dots n]$ and $Y[1 \dots m]$.

if $(0 == n \text{ or } 0 == m)$ return 0

if $(X[n] == Y[m])$

return $1 + LCS(n-1, m-1)$

return $\max(LCS(n, m-1),$
 $LCS(n-1, m))$

LCS: Iterative Solution*

LCS(n,m): // recursive for reference

if $0 == n$ or $0 == m$ then

return 0

else if $X[n] = Y[m]$ then

return $1 + \text{LCS}(n-1, m-1)$

else

return $\max(\text{LCS}(n-1, m), \text{LCS}(n, m-1))$

for ($j=0 \dots m$)
LCS[0, j] = 0

for $i=1 \dots n$
LCS[i, 0] = 0

declare LCS[0..n, 0..m]

*

// non-base:

for $i=1 \dots n$

for $j=1 \dots m$

// fill in LCS[i, j] as per

if $X[i] == Y[j]$

LCS[i, j] = 1 +
LCS[i-1, j-1]

$O(mn)$

and other case too! above

LCS Example:

		M	O	R	N	I	N	G
	0	0	0	0	0	0	0	0
T	0							
R	0							
I	0							
A	0							
N	0							
G	0							
L	0							
E	0							

good
reinforcement

Finding the LCS itself

► We have $LCS[]$ filled in.

$i \leftarrow n, j \leftarrow m, S \leftarrow \text{empty stack}$

while **do**

// is $x[i]$ and $y[j]$ part of output?

good
reinforcement

Did we do everything we need to do?

- ▶ Describe *in English* the function
 - ▶ Not *how* it works (yet)
 - ▶ Yes *what it solves*.
 - ▶ Skipping this step = 0 on problem
- ▶ Give that function a meaningful variable name.
 - ▶ Not “OPT” or “DP” or “table”
 - ▶ Not a single letter either.
- ▶ Give recursive formulation.
- ▶ Describe the iterative running time.
 - ▶ Was cut-off in Friday's slides. :(