IP - The Internet Protocol

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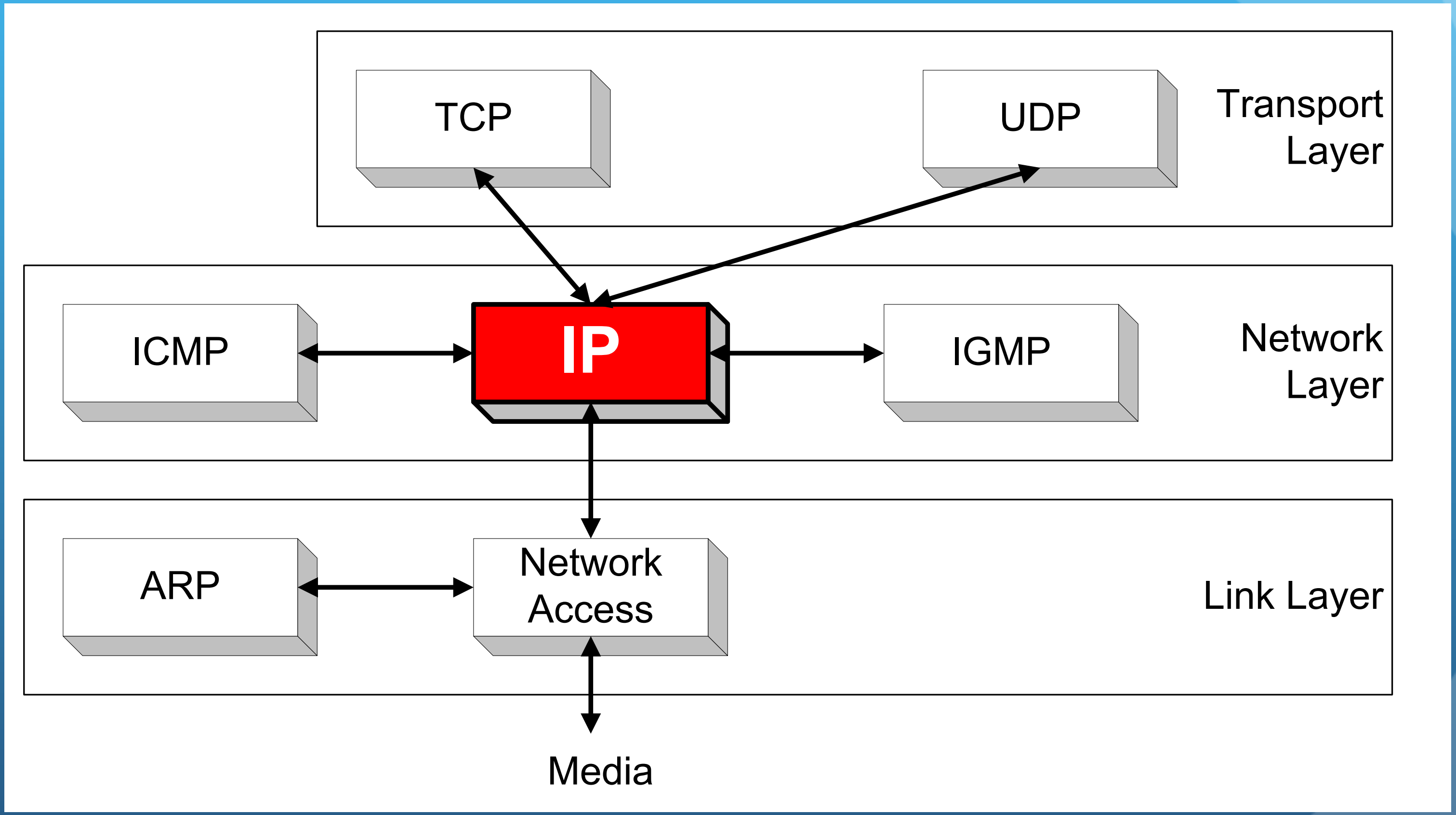
Overview

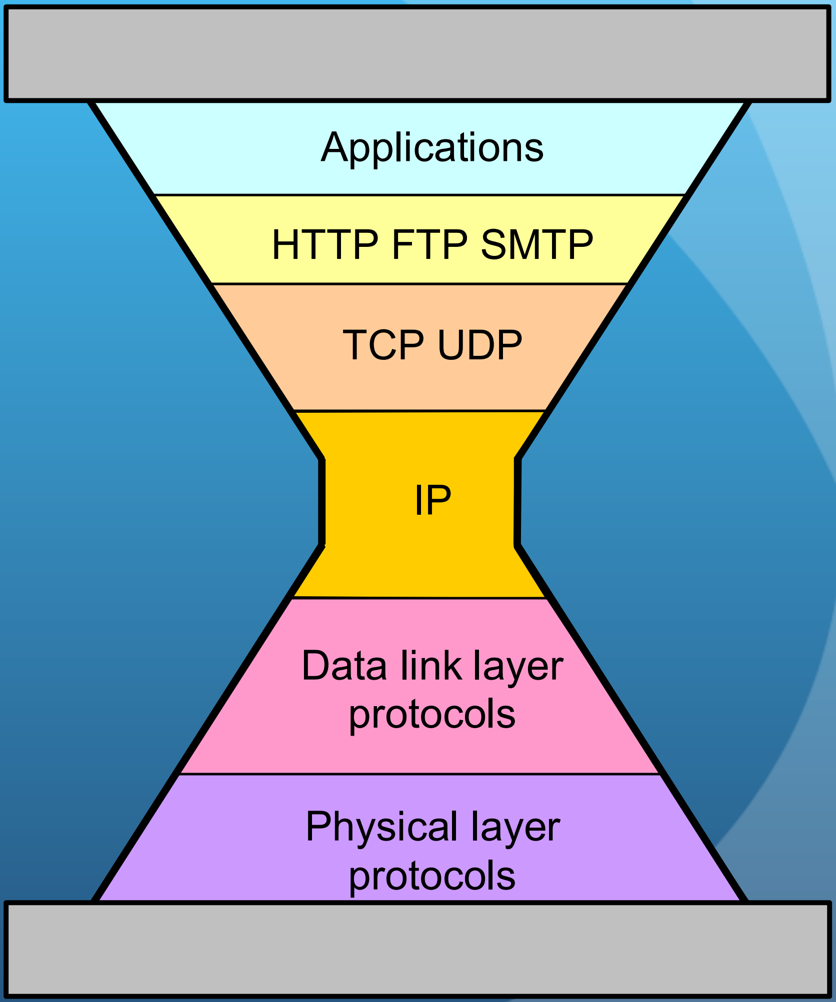
- IP (Internet Protocol) is a Network Layer Protocol.
- Several versions – most popular (IPv4). It is specified in RFC 891.
- Gaining popularity is IPv6 due to increased addressing space and security handling.

IP and the Internet Architecture



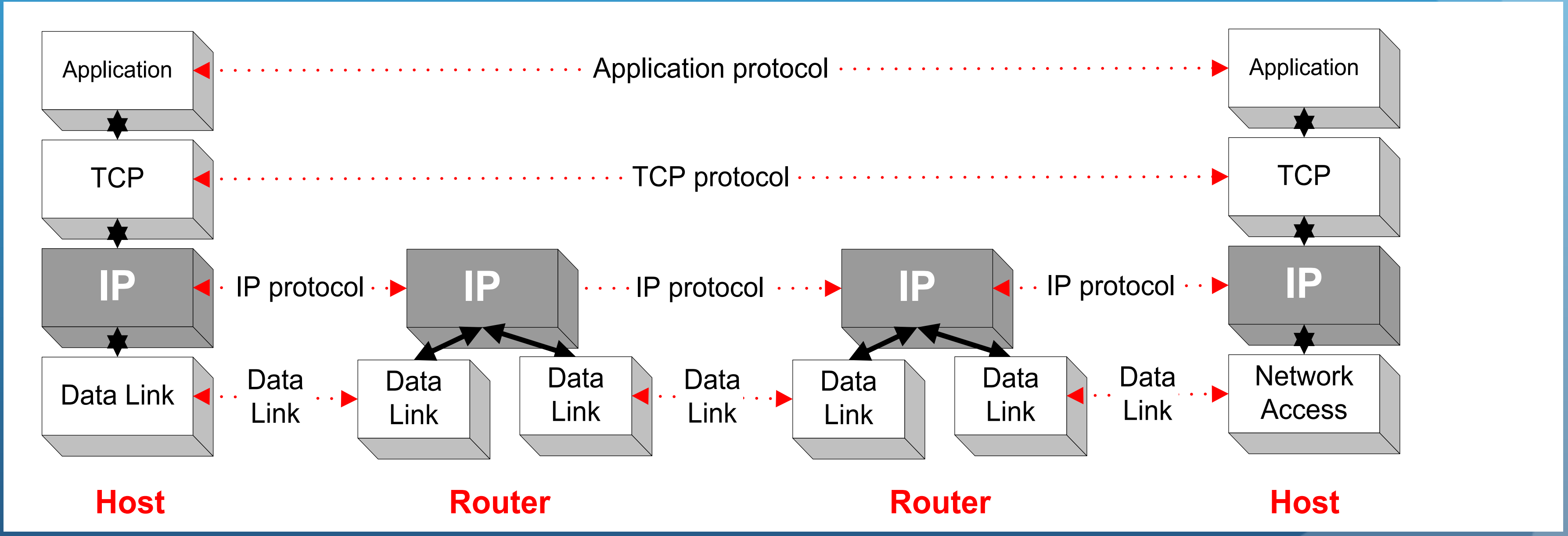
IP: The waist of the hourglass

- **IP is the waist of the hourglass of the Internet protocol architecture**
- Multiple higher-layer protocols
- Multiple lower-layer protocols
- Only one protocol at the network layer.



The Internet protocol

- IP is the highest layer protocol which is implemented at BOTH routers and hosts

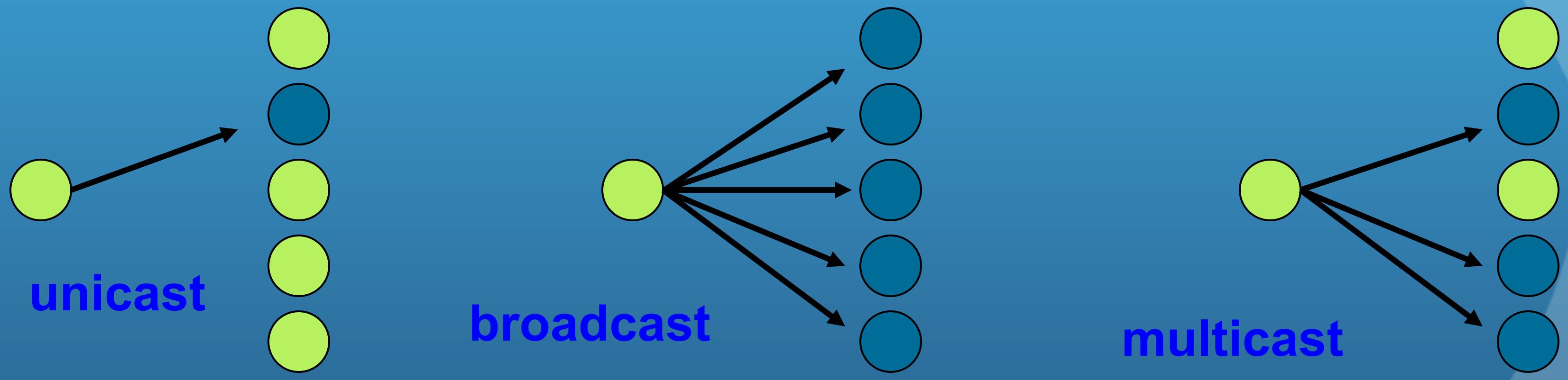


IP Service

- Delivery service of IP is minimal
- IP provide provides an unreliable connectionless best effort service (also called: "datagram service").
  - **Unreliable:** IP does not make an attempt to recover lost packets
  - **Connectionless:** Each packet ("datagram") is handled independently. IP is not aware that packets between hosts may be sent in a logical sequence
  - **Best effort:** IP does not make guarantees on the service (no throughput guarantee, no delay guarantee,…)
- Consequences:
  - Higher layer protocols have to deal with losses or with duplicate packets
  - Packets may be delivered out-of-sequence

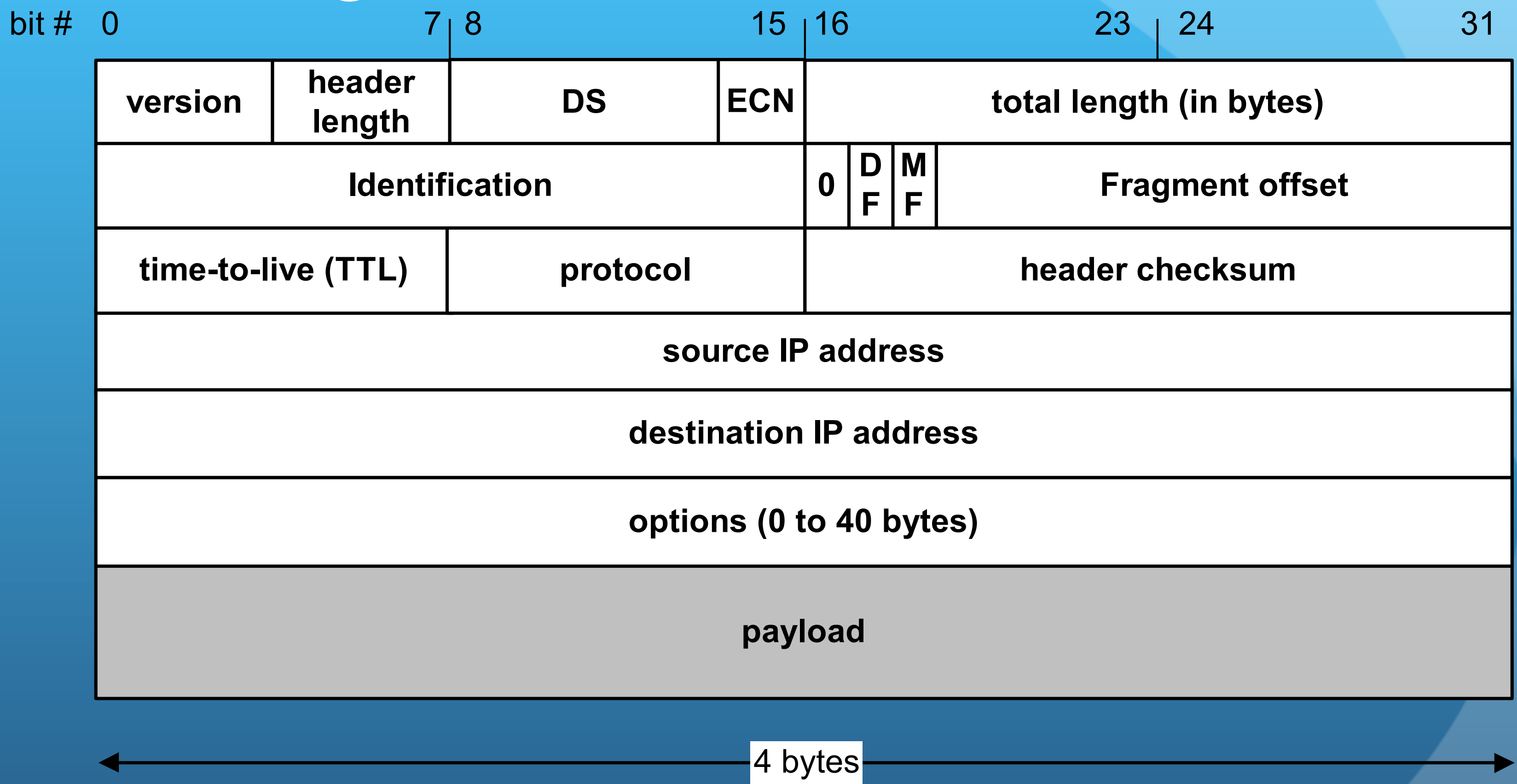
IP Service

- IP supports the following services:
  - one-to-one (unicast)
  - one-to-all (broadcast)
  - one-to-several (multicast)



- IP multicast also supports a many-to-many service.
- IP multicast requires support of other protocols (IGMP, multicast routing)

IP Datagram Format



- 20 bytes ≤ Header Size < 2^4 x 4 bytes = 60 bytes
- 20 bytes ≤ Total Length < 2^{16} bytes = 65536 bytes

IP Datagram Format

- **Question:** In which order are the bytes of an IP datagram transmitted?

- **Answer:**
  - Transmission is row by row
  - For each row:
    1. First transmit bits 0-7
    2. Then transmit bits 8-15
    3. Then transmit bits 16-23
    4. Then transmit bits 24-31

- This is called **network byte** order or **big endian** byte ordering.

- **Note:** some computers store 32-bit words in little endian format.

Fields of the IP Header

- **Version (4 bits):** current version is 4, next version will be 6.
- **Header length (4 bits):** length of IP header, in multiples of 4 bytes
- **DS/ECN field (1 byte)**
  - This field was previously called as Type-of-Service (TOS) field. The role of this field has been re-defined, but is "backwards compatible" to TOS interpretation
  - Differentiated Service (DS) (6 bits):
    - Used to specify service level (currently not supported in the Internet)
  - Explicit Congestion Notification (ECN) (2 bits):
    - Feedback mechanism used by TCP

Fields of the IP Header

- **Identification (16 bits):** Unique identification of a datagram from a host. Incremented whenever a datagram is transmitted

- **Flags (3 bits):**
  - First bit always set to 0
  - DF bit (Do not fragment)
  - MF bit (More fragments)

  Will be explained later→ Fragmentation

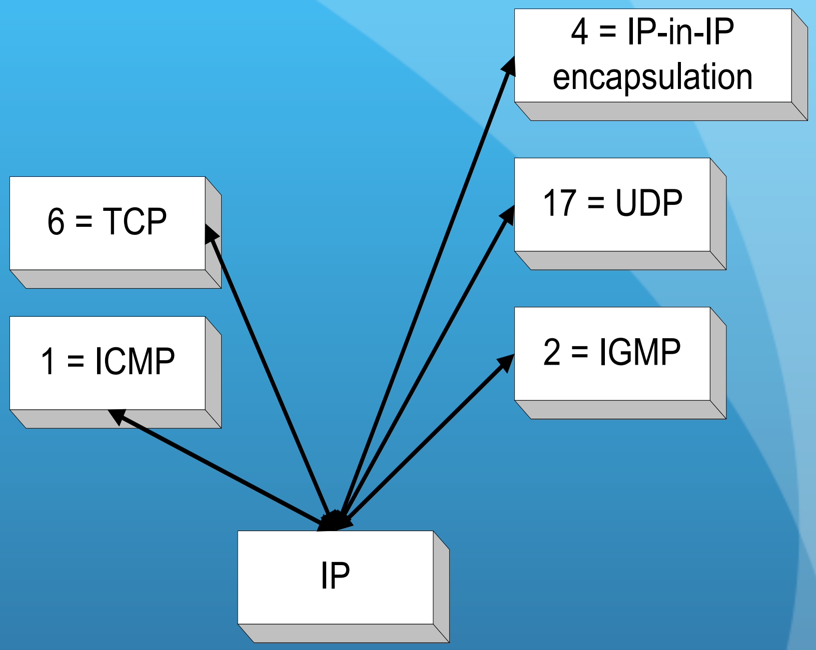
Fields of the IP Header

- **Time To Live (TTL) (1 byte):**
  - Specifies longest paths before datagram is dropped
  - Role of TTL field: Ensure that packet is eventually dropped when a routing loop occurs

  Used as follows:
  - Sender sets the value (e.g., 16)
  - Each router decrements the value by 1
  - When the value reaches 0, the datagram is dropped

Fields of the IP Header

- **Protocol (1 byte):**
  - Specifies the higher-layer protocol.
  - Used for demultiplexing to higher layers.
- **Header checksum (2 bytes):** A simple 16-bit long checksum which is computed for the header of the datagram.



Fields of the IP Header

- **Options:**
  - Security restrictions
  - Record Route: each router that processes the packet adds its IP address to the header.
  - Timestamp: each router that processes the packet adds its IP address and time to the header.
  - (loose) Source Routing: specifies a list of routers that must be traversed.
  - (strict) Source Routing: specifies a list of the only routers that can be traversed.
- **Padding:** Padding bytes are added to ensure that header ends on a 4-byte boundary

Maximum Transmission Unit

- Maximum size of IP datagram is 65535, but the data link layer protocol generally imposes a limit that is much smaller
- For example:
  - Ethernet frames have a maximum payload of 1500 bytes→ IP datagrams encapsulated in Ethernet frame cannot be longer than 1500 bytes
- The limit on the maximum IP datagram size, imposed by the data link protocol is called **maximum transmission unit (MTU)**
- MTUs for various data link layers:

Ethernet:	1500	FDDI:	4352
802.3:	1492	ATM AAL5:	9180
802.5:	4464	PPP:	296

- What if the size of an IP datagram exceeds the MTU?

  IP datagram is fragmented into smaller units.
- What if the route contains networks with different MTUs?

Maximum Transmission Unit

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- Example:
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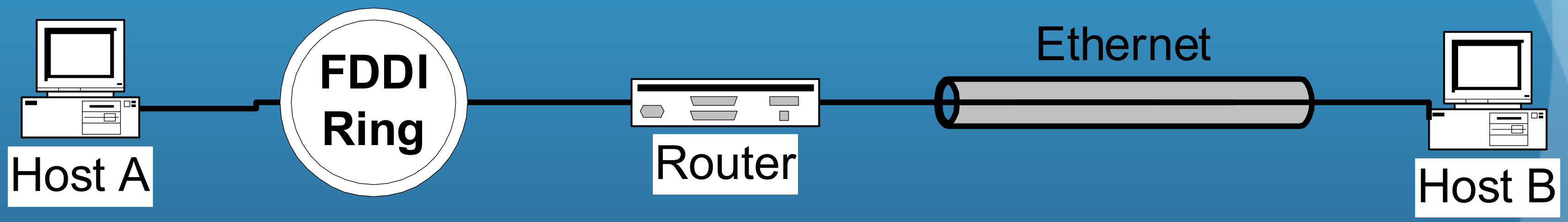
    → IP datagrams encapsulated in Ethernet frame cannot be longer than 1500 bytes
- The limit on the maximum IP datagram size, imposed by the data link protocol is called **maximum transmission unit (MTU)**
- MTUs for various data link protocols:

Ethernet:	1500	FDDI:	4352
802.3:	1492	ATM AAL5:	9180
802.5:	4464	PPP:	negotiated

IP Fragmentation

- What if the size of an IP datagram exceeds the MTU?
- IP datagram is fragmented into smaller units.

- What if the route contains networks with different MTUs?
- IP datagram is fragmented into smaller and smaller units

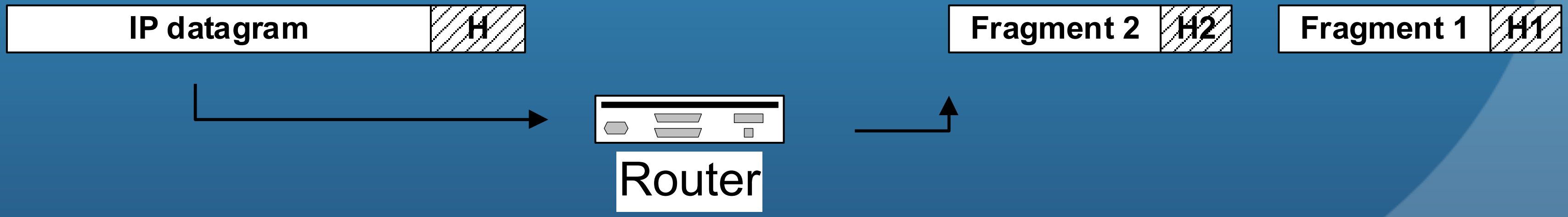


MTUs: FDDI: 4352 Ethernet: 1500

- **Fragmentation**:
  - IP router splits the datagram into several datagrams
  - Fragments are reassembled at receiver

Where is Fragmentation done?

- Fragmentation can be done at the sender or at intermediate routers
- The same datagram can be fragmented several times.
- Reassembly of original datagram is only done at destination hosts !!



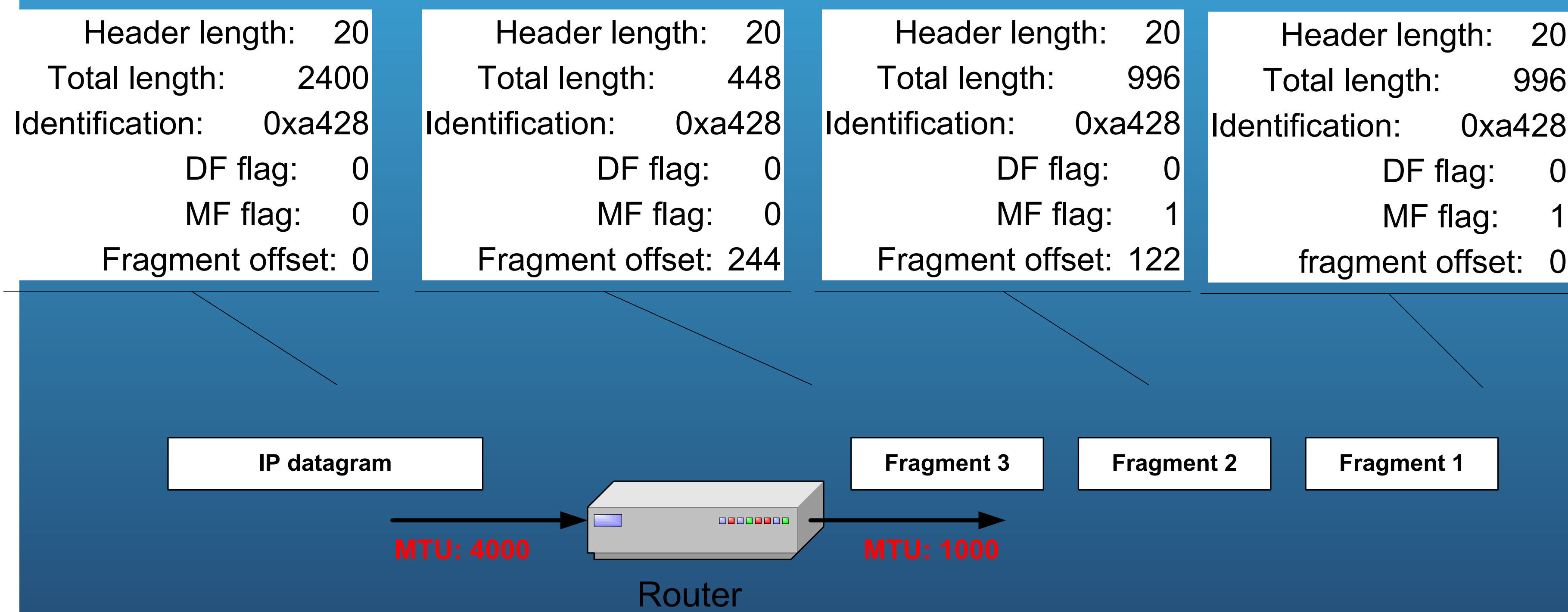
What is involved in Fragmentation?

version	header length	DS	ECN	total length (in bytes)			
Identification				0	DF	MF	Fragment offset
time-to-live (TTL)		protocol		header checksum			

Identification	When a datagram is fragmented, the identification is the same for all fragments
Flags	
DF bit is set:	Datagram cannot be fragmented and must be discarded if MTU is too small
MF bit set:	This datagram is part of a fragment and an additional fragment follows this one
Fragment offset	Offset of the payload of the current fragment in the original datagram
Total length	Total length of the current fragment

Example of Fragmentation

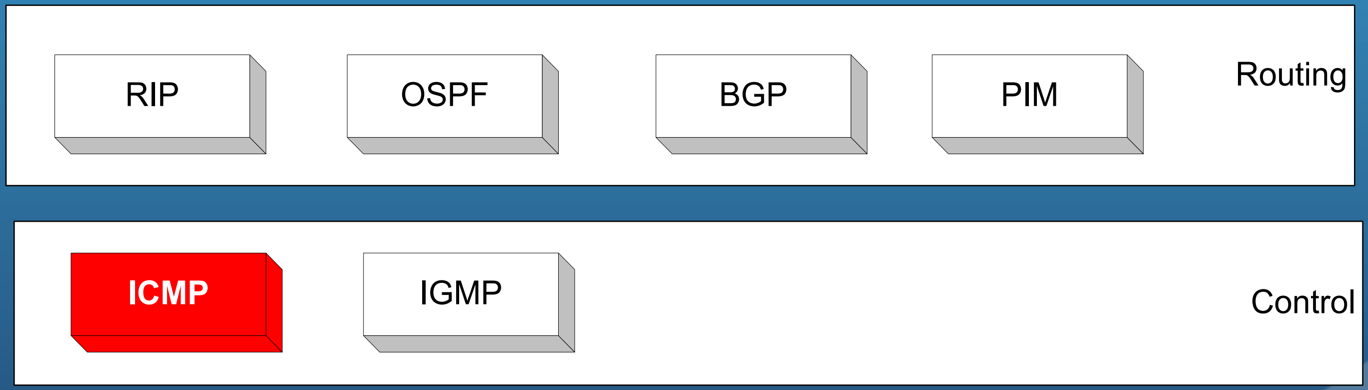
- A datagram with size 2400 bytes must be fragmented according to an MTU limit of 1000 bytes



Internet Control Message Protocol (ICMP)

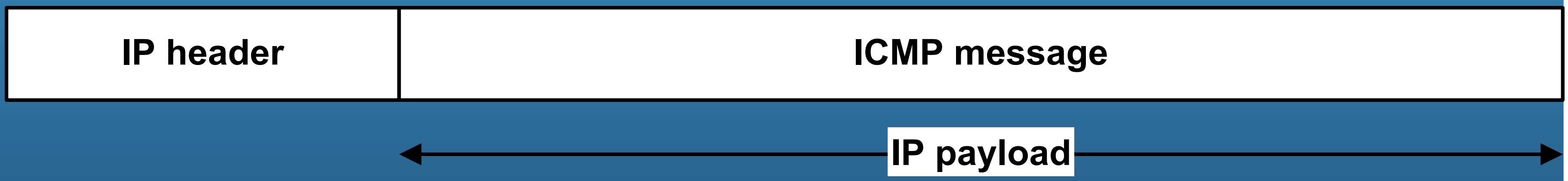
Overview

- The IP (Internet Protocol) relies on several other protocols to perform necessary control and routing functions:
  - Control functions (ICMP)
  - Multicast signaling (IGMP)
  - Setting up routing tables (RIP, OSPF, BGP, PIM, …)

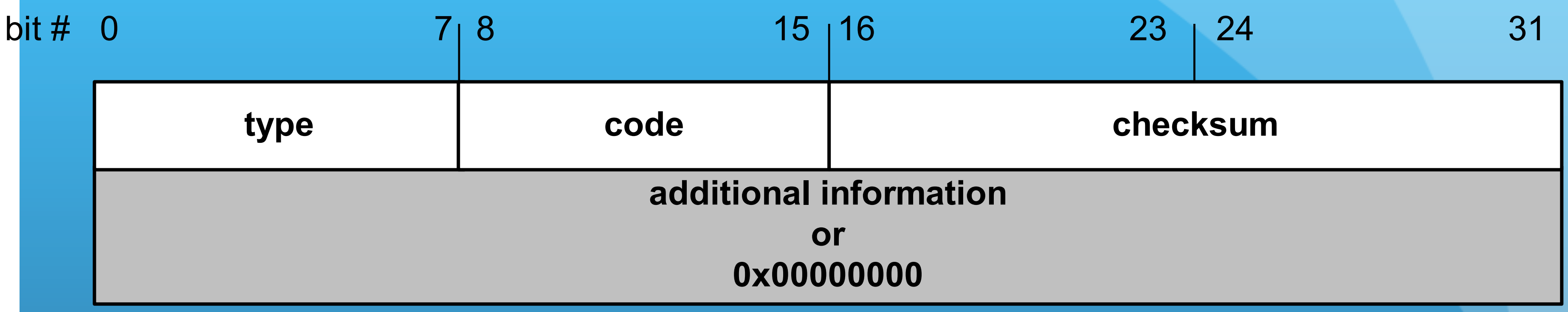


ICMP

- The **Internet Control Message Protocol (ICMP)** is a helper protocol that supports IP with:
  - Error reporting
  - Simple queries
- ICMP messages are encapsulated as IP datagrams:



ICMP message format



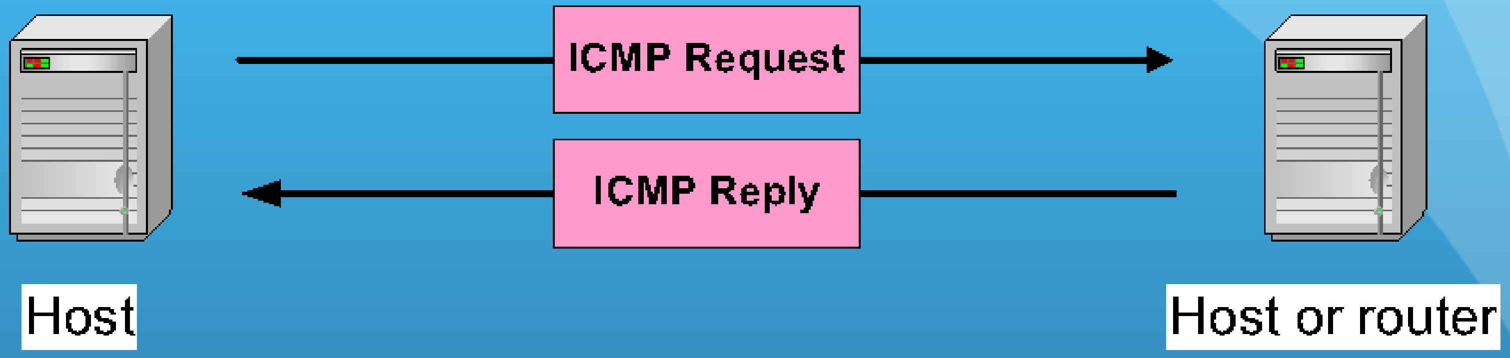
4 byte header:

- Type (1 byte): type of ICMP message
- Code (1 byte): subtype of ICMP message
- Checksum (2 bytes): similar to IP header checksum. Checksum is calculated over entire ICMP message

If there is no additional data - 4 bytes set to zero

→ each ICMP messages is at least 8 bytes long

ICMP Query message



ICMP query:

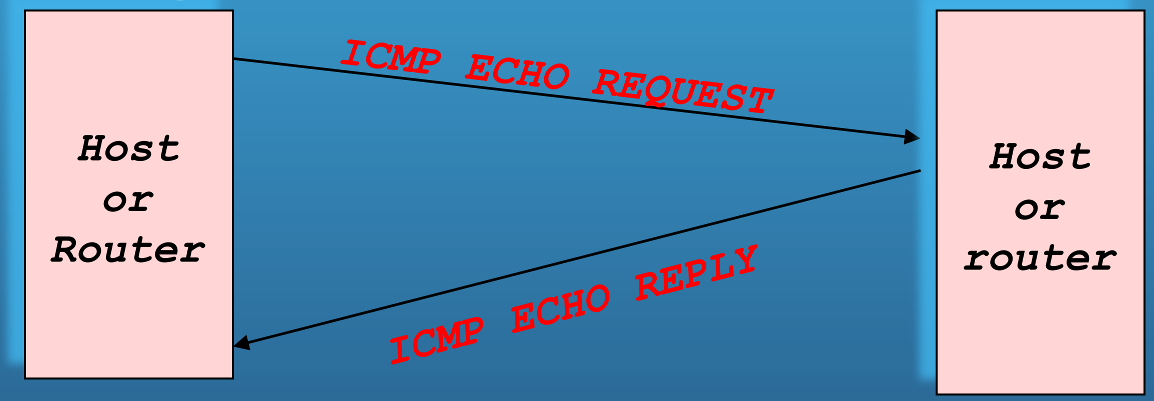
- Request sent by host to a router or host
- Reply sent back to querying host

Examples of ICMP Queries

Type/Code:	Description	
8/0	Echo Request	The ping command uses Echo Request/ Echo Reply
0/0	Echo Reply	
13/0	Timestamp Request	
14/0	Timestamp Reply	
10/0	Router Solicitation	
9/0	Router Advertisement	

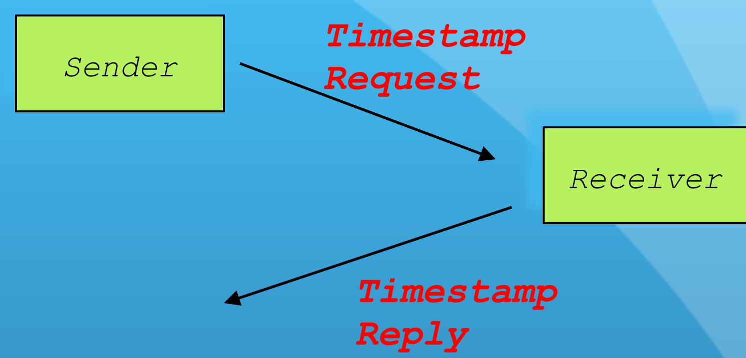
Example of a Query: Echo Request and Reply

- Ping’s are handled directly by the kernel
- Each Ping is translated into an ICMP Echo Request
- The Ping’ed host responds with an ICMP Echo Reply



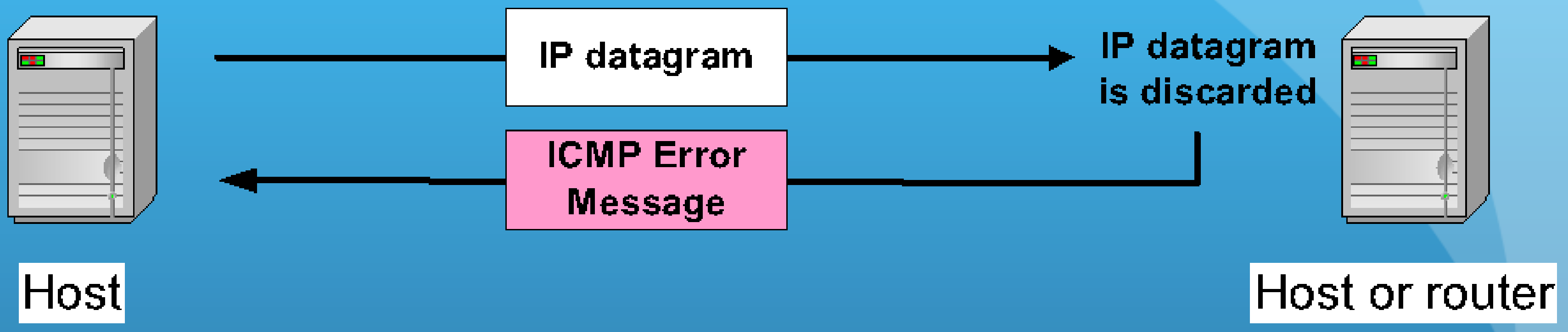
Example of a Query: ICMP Timestamp

- A system (host or router) asks another system for the current time.
- Time is measured in milliseconds after midnight UTC (Universal Coordinated Time) of the current day
- Sender sends a request, receiver responds with reply



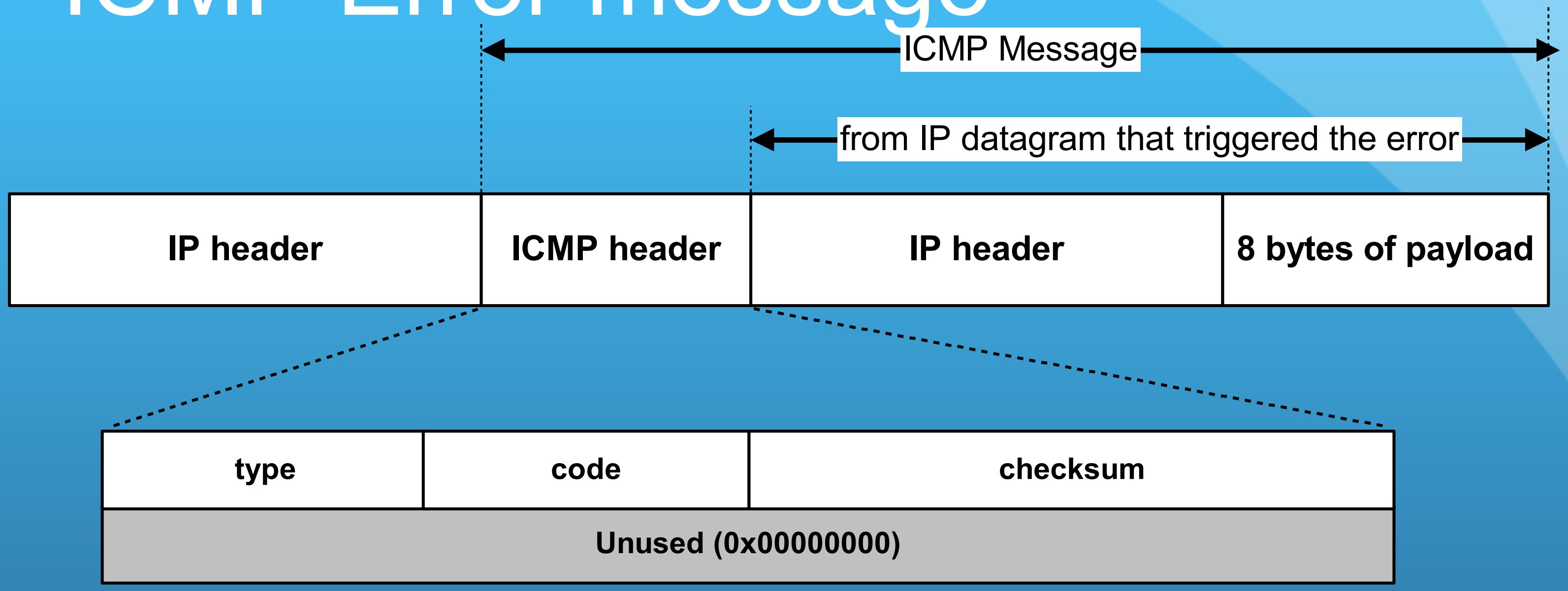
Type (= 17 or 18)	Code (=0)	Checksum
identifier		sequence number
32-bit sender timestamp		
32-bit receive timestamp		
32-bit transmit timestamp		

ICMP Error message



- **ICMP error messages report error conditions**
- **Typically sent when a datagram is discarded**
- **Error message is often passed from ICMP to the application program**

ICMP Error message



- ICMP error messages include the complete IP header and the first 8 bytes of the payload (typically: UDP, TCP)

Common ICMP Error messages

Type	Code	Description	
3	0–15	Destination unreachable	Notification that an IP datagram could not be forwarded and was dropped. The code field contains an explanation.
5	0–3	Redirect	Informs about an alternative route for the datagram and should result in a routing table update. The code field explains the reason for the route change.
11	0, 1	Time exceeded	Sent when the TTL field has reached zero (Code 0) or when there is a timeout for the reassembly of segments (Code 1)
12	0, 1	Parameter problem	Sent when the IP header is invalid (Code 0) or when an IP header option is missing (Code 1)

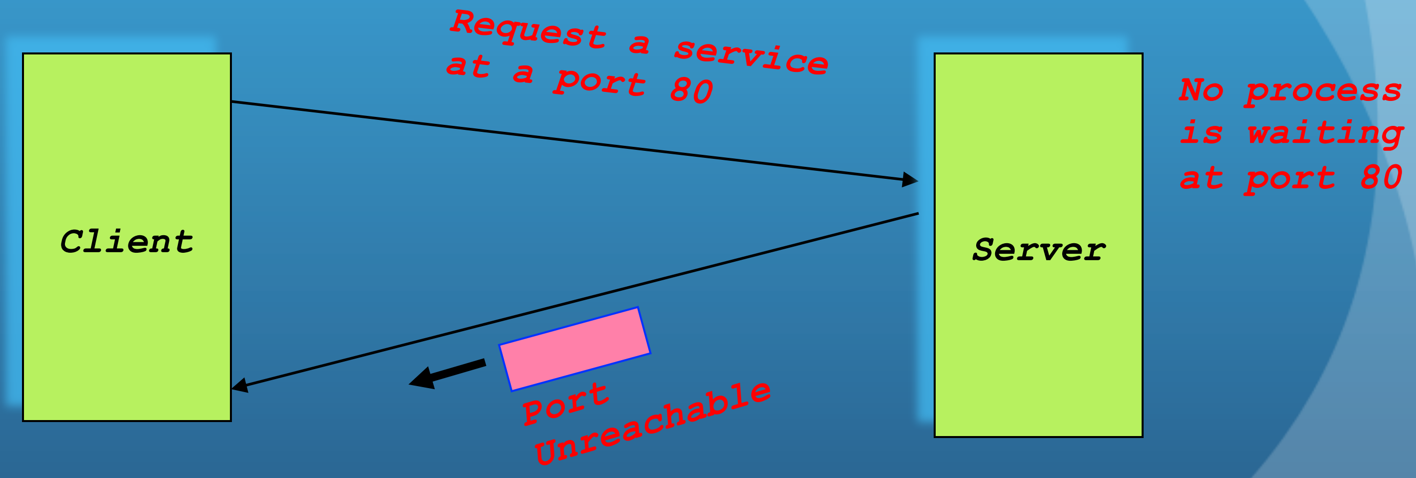
Some subtypes of the “Destination Unreachable”

Code	Description	Reason for Sending
0	Network Unreachable	No routing table entry is available for the destination network.
1	Host Unreachable	Destination host should be directly reachable, but does not respond to ARP Requests.
2	Protocol Unreachable	The protocol in the protocol field of the IP header is not supported at the destination.
3	Port Unreachable	The transport protocol at the destination host cannot pass the datagram to an application.
4	Fragmentation Needed and DF Bit Set	IP datagram must be fragmented, but the DF bit in the IP header is set.

Example: ICMP Port Unreachable

- RFC 792: If, in the destination host, the IP module cannot deliver the datagram because the indicated protocol module or process port is not active, the destination host may send a destination unreachable message to the source host.

- Scenario:



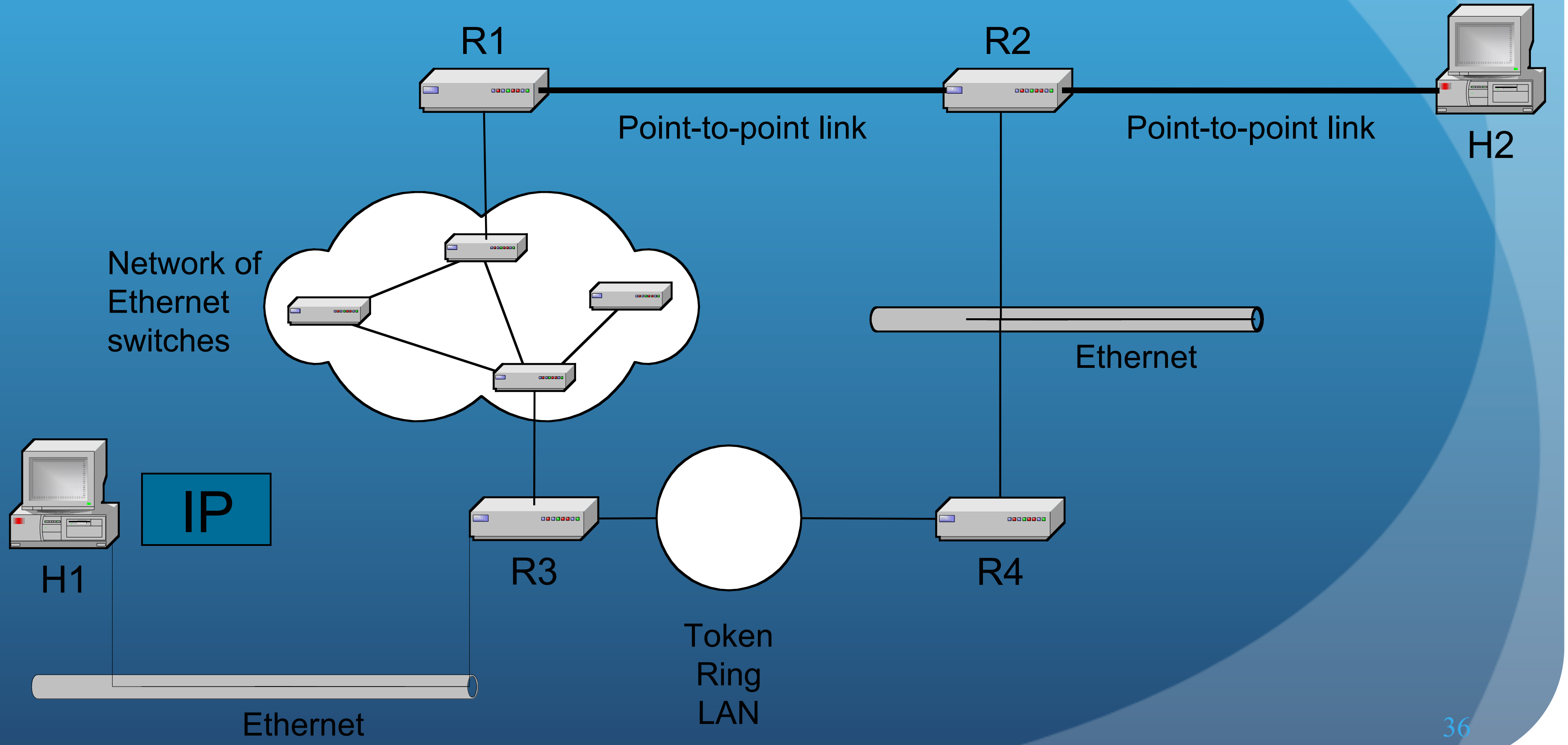
IP Forwarding

Overview

- Internet is a collection of networks
- IP provides an end-to-end delivery service between hosts
- The delivery service is realized with the help of IP routers

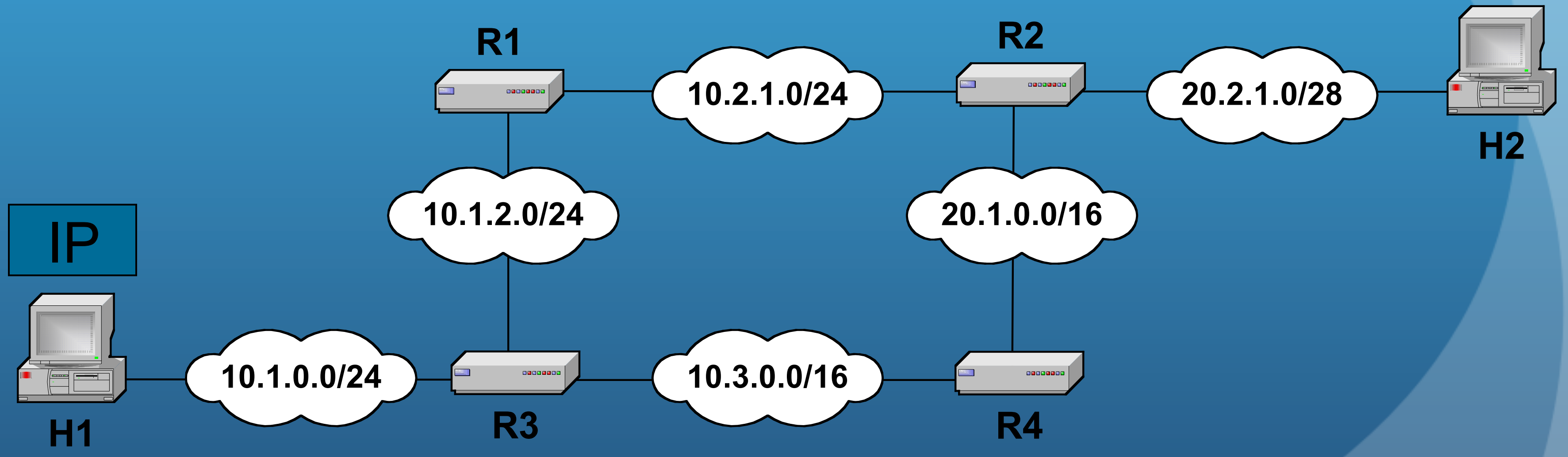
Delivery of an IP datagram

- View at the data link layer layer:
  - Internetwork is a collection of LANs or point-to-point links or switched networks that are connected by routers



Delivery of an IP datagram

- View at the IP layer:
  - An IP network is a logical entity with a network number
  - We represent an IP network as a "cloud"
  - The IP delivery service takes the view of clouds, and ignores the data link layer view



Tenets of end-to-end delivery of datagrams

- The following conditions must hold so that an IP datagram can be successfully delivered

1. The network prefix of an IP destination address must correspond to a unique data link layer network (=LAN or point-to-point link or switched network).
2. Routers and hosts that have a common network prefix must be able to exchange IP datagrams using a data link protocol (e.g., Ethernet, PPP)
3. An IP network is formed when a data link layer network is connected to at least one other data link layer network via a router.

Routing tables

- Each router and each host keeps a **routing table** which tells the router how to process an outgoing packet
- Main columns:
  1. **Destination address:** where is the IP datagram going to?
  2. **Next hop or interface:** how to send the IP datagram?
- Routing tables are set so that a datagram gets closer to the its destination

Routing table of a host or router

IP datagrams can either be directly delivered ("direct") or are sent to a next hop router ("R4")

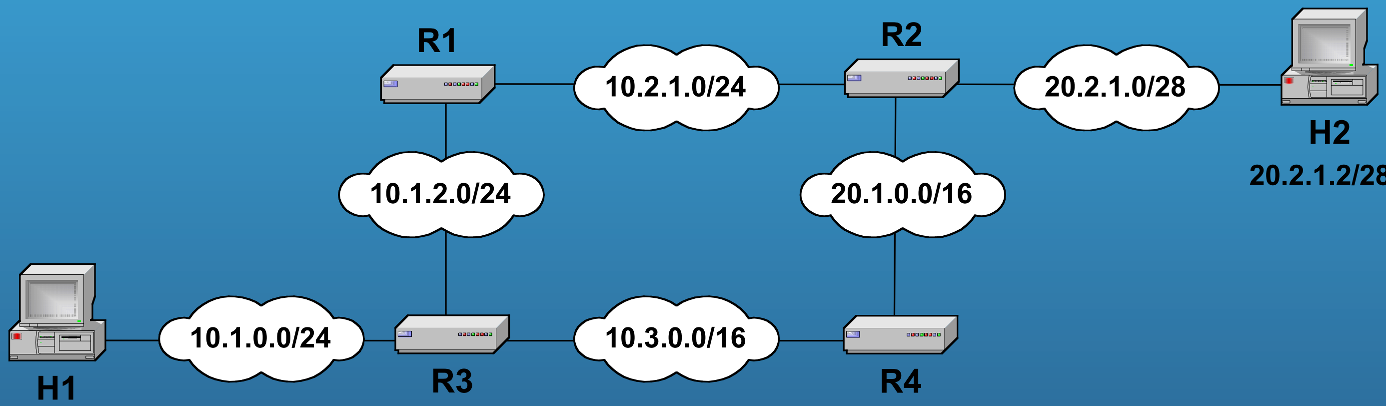
Destination	**Next Hop**
20.2.1.0/28	R4
10.1.0.0/24	direct
10.1.2.0/24	direct
10.2.1.0/24	R4
10.3.1.0/24	direct
20.1.0.0/16	R4

Delivery with routing tables

Destination	Next Hop
10.1.0.0/24	R3
10.1.2.0/24	direct
10.2.1.0/24	direct
10.3.1.0/24	R3
20.2.0.0/16	R2
30.1.1.0/28	R2

Destination	Next Hop
10.1.0.0/24	R1
10.1.2.0/24	R1
10.2.1.0/24	direct
10.3.1.0/24	R4
20.1.0.0/16	direct
20.2.1.0/28	direct

Destination	Next Hop
10.1.0.0/24	R2
10.1.2.0/24	R2
10.2.1.0/24	R2
10.3.1.0/24	R2
20.1.0.0/16	R2
20.2.1.0/28	direct



to:
20.2.1.2

Destination	Next Hop
10.1.0.0/24	direct
10.1.2.0/24	R3
10.2.1.0/24	R3
10.3.1.0/24	R3
20.1.0.0/16	R3
20.2.1.0/28	R3

Destination	Next Hop
10.1.0.0/24	direct
10.1.2.0/24	direct
10.2.1.0/24	R4
10.3.1.0/24	direct
20.1.0.0/16	R4
20.2.1.0/28	R4

Destination	Next Hop
10.1.0.0/24	R3
10.1.2.0/24	R3
10.2.1.0/24	R2
10.3.1.0/24	direct
20.1.0.0/16	direct
20.2.1.0/28	R2

Delivery of IP datagrams

- There are two distinct processes to delivering IP datagrams:

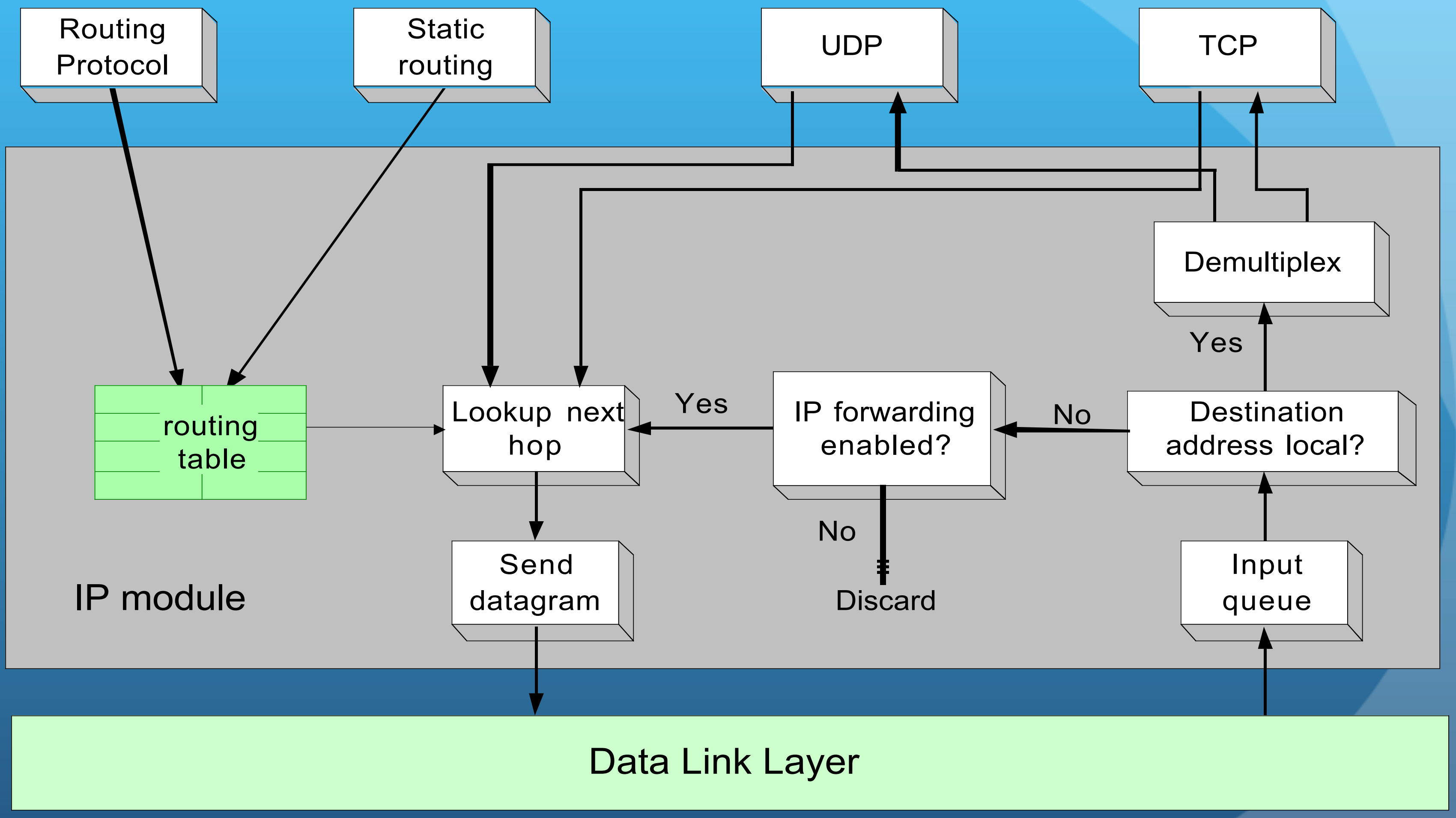
  1. **Forwarding:** How to pass a packet from an input interface to the output interface?

  2. **Routing:** How to find and setup the routing tables?

- Forwarding must be done as fast as possible:
  - on routers, is often done with support of hardware
  - on PCs, is done in kernel of the operating system

- Routing is less time-critical
  - On a PC, routing is done as a background process

Processing of an IP datagram in IP



Processing of an IP datagram in IP

- Processing of IP datagrams is very similar on an IP router and a host
- **Main difference:**
  "IP forwarding" is enabled on router and disabled on host
- **IP forwarding enabled**
  → if a datagram is received, but it is not for the local system, the datagram will be sent to a different system
- **IP forwarding disabled**
  → if a datagram is received, but it is not for the local system, the datagram will be discarded

Processing of an IP datagram at a router

Receive an IP datagram →

1. IP header validation
2. Process options in IP header
3. Parsing the destination IP address
4. Routing table lookup
5. Decrement TTL
6. Perform fragmentation (if necessary)
7. Calculate checksum
8. Transmit to next hop
9. Send ICMP packet (if necessary)

Routing table lookup

- When a router or host need to transmit an IP datagram, it performs a routing table lookup
- **Routing table lookup:** Use the IP destination address as a key to search the routing table.
- Result of the lookup is the IP address of a next hop router, or the name of a network interface

Destination address	**Next hop**
network prefix *or* host IP address *or* loopback address *or* default route	IP address of next hop router* *or* Name of a network interface

*__Note__: A router has many IP addresses. The IP address in the routing table refers to the address of the network interface on the same directly connected network.

Type of routing table entries

- **Network route**
  - Destination addresses is a network address (e.g., 10.0.2.0/24)
  - Most entries are network routes
- **Host route**
  - Destination address is an interface address (e.g., 10.0.1.2/32)
  - Used to specify a separate route for certain hosts
- **Default route**
  - Used when no network or host route matches
  - The router that is listed as the next hop of the default route is the **default gateway (for Cisco: "gateway of last resort)**
- **Loopback address**
  - Routing table for the loopback address (127.0.0.1)
  - The next hop lists the loopback (lo0) interface as outgoing interface

Longest Prefix Match

- **Longest Prefix Match:** Search for the routing table entry that has the longest match with the prefix of the destination IP address

1. Search for a match on all 32 bits
2. Search for a match for 31 bits

   …..

32. Search for a match on 0 bits

Host route, loopback entry
→ 32-bit prefix match

Default route is represented as 0.0.0.0/0
→ 0-bit prefix match

128.143.71.21

Destination address	Next hop
10.0.0.0/8	R1
128.143.0.0/16	R2
128.143.64.0/20	R3
128.143.192.0/20	R3
128.143.71.0/24	R4
128.143.71.55/32	R3
default	R5

The longest prefix match for 128.143.71.21 is entry 128.143.71.0/24 with 24 bit match
-> Datagram sent to R4

Route Aggregation

- Longest prefix match algorithm permits the aggregation of prefixes with identical next hop address to a single entry
- This contributes significantly to reducing the size of routing tables of Internet routers

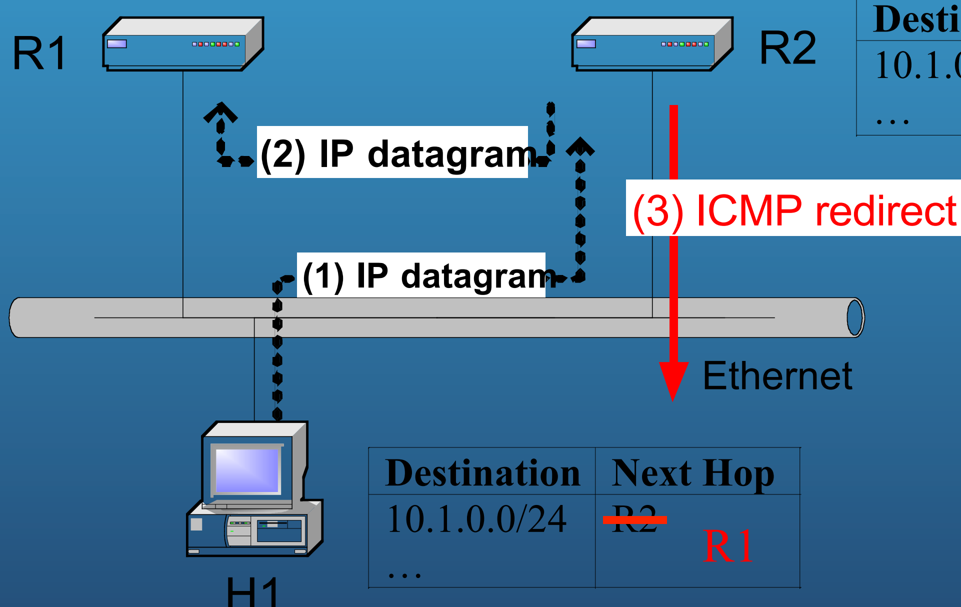
Destination	Next Hop
10.1.0.0/24	R3
10.1.2.0/24	direct
10.2.1.0/24	direct
10.3.1.0/24	R3
20.2.0.0/16	R2
20.1.1.0/28	R2

Destination	Next Hop
10.1.0.0/24	R3
10.1.2.0/24	direct
10.2.1.0/24	direct
10.3.1.0/24	R3
20.0.0.0/14	R2

Routing table manipulations with ICMP

- When a router detects that an IP datagram should have gone to a different router, the router (here R2)
  - forwards the IP datagram to the correct router
  - sends an ICMP redirect message to the host

- Host uses ICMP message to update its routing table



Destination	Next Hop
10.1.0.0/24	R1
…	

Destination	Next Hop
10.1.0.0/24	~~R2~~ R1
…	

ICMP Router Solicitation ICMP Router Advertisement

- After bootstrapping a router broadcasts an **ICMP router solicitation.**
- In response, routers send an **ICMP router advertisement** message
- Also, routers periodically broadcast **ICMP router advertisement**

This is sometimes called the Router Discovery Protocol