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# **Robust Localization Protocols and Algorithms in Wireless Sensor Networks Using UWB**

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Localization has many important applications in wireless sensor networks (WSNs). A variety of technologies, such as acoustic, infrared, and UWB (ultra-wide band) media have been utilized for localization purposes. In this paper, we propose a holistic, bottom-up design of a UWB-based communication architecture and related protocols for localization in WSNs. A new UWB coding method, called U-BOTH (UWB based on Orthogonal Variable Spreading Factor and Time Hopping), is utilized for minimum interference communication, and an ALOHA-type channel access method and a message exchange protocol are used to collect sensor location and signal strength information in WSNs. After establishing the UWB path loss model, we apply the maximum likelihood estimation (MLE) method to compute the distances between neighbor nodes using the RSS information. Then, we propose NMDS-MLE (Non-metric Multidimensional Scaling and Maximum Likelihood Estimation) localization algorithms by matching the sensor coordinate estimates with the distance estimates derived from the path loss model. The performance of the system is validated using theoretic analysis, simulation and real testbed experiments.

*Keywords:* Orthogonal Variable Spreading Factor (OVSF), Time Hopping (TH), Ultra-Wide Band (UWB), localization, ranging.

## **1 INTRODUCTION**

Large-scale economic wireless sensor networks (WSNs) are widely deployed for environmental monitoring and control operations. Object tracking and

localization are two important capabilities in many WSN applications [3]. The basic approach to a WSN localization is to infer distances to anchors (coordinate-known nodes), then to derive targets (coordinate-unknown nodes) by trilateration or other estimation algorithms. The first step is called "ranging", and the second step is called "localization".

So far, various ranging solutions have been proposed based on two major ranging techniques: 1) time of arrival (ToA) [16], time difference of arrival (TDOA) and angle of arrival (AOA), as used by GPS systems, 2) path loss models describing the propagation characteristics of radio received signal strength (RSS) [4] or acoustic signal strength [21], 3) Range-free techniques using hop count or centroid methods [9]. In this paper, we adopt the path loss model to derive ranges, because it only requires signal strength information, which is easy to collect in WSNs, in contrast to expensive synchronization or topology information collection used in other techniques [8].

Ranging algorithms based on path loss model depend on the wireless medium and signal transmission methods. In order to provide precision ranging, we utilize the UWB (ultra-wide band) transmission and coding technologies in both indoor and outdoor environments. Beside providing high data bandwidth, UWB exhibits excellent resistance to co-channel interference. IEEE 802.15.4a has appeared as the *de facto* standard to provide low-power long-distant low-data-rate service for real-time communication and precise ranging and localization applications [1, 14].

Of the different UWB transmission techniques, Impulse Radio Ultra Wideband (IR-UWB) is most attractive for localization purposes in WSNs [22]. However, existing coding algorithms for IR-UWB communication systems, such as DS-UWB (Direct Sequence UWB) and TH-UWB (Time Hopping UWB) [13] have failed to guarantee high quality localization due to multipath and multi-user interference. In this paper, we apply the Orthogonal Variable Spread Factor (OVSF) coding algorithm in IR-UWB networks to solve the multi-user interference problem in data transmissions.

Once the approximate distances between a target and a subset of anchor points are derived in the WSNs, the target coordinate can be derive by localization algorithms. Savarese *et al.* presented a trilateration algorithm based on least squares (LS) method in large-scale WSNs [18]. Caplun *et al.* [7] proposed a GPS-free positioning system for mobile ad hoc networks, by first establishing the local coordinates of two-hop neighbors with each node as the origin, then tuning these local coordinates to the global coordinates of the entire system. DV-coordinate (distance vector) algorithm [15] applied a similar idea.

Different from trilateration algorithms, the MDS (Multidimensional Scaling) method translates data objects represented in *N*-dimensional similarity metric to locations in the 2- or 3-dimensional Euclidean space [5]. The relative map provides partial and relative inter-nodal relationships, whereas the absolute map conforms to the relationships in the relative map. Therefore,

MDS requires less information and configuration overhead than other localization algorithms in WSNs, and provides strong resilience to measurement errors.

Several variants of MDS-based localization algorithms exist. MDS-MAP uses connectivity information (whether or not two devices are in range) for localization [20]. MDS-MAP(P) improved the basic MDS-MAP on anisotropic topologies [19] by building the local relative map of a small subnetwork for each node using MDS, then patching them to form a global relative map. However, most of MDS algorithms were based on the assumption that the proximity between nodes in the relative map should be proportional to Euclidean distances in the absolute map in the underlying quantitative transformation function, which is not flexible or robust.

In this paper, we present the NMDS-MLE (Non-metric MDS and Maximum Likelihood Estimation) localization algorithm. In contrast to MDS which relies on accurate distance measurements to derive the Euclidean coordinates, NMDS only needs pair-wise comparative distance relations, represented with "less than", to adjust the Euclidean coordinates.

In NMDS-MLE, we deploy a small number of anchor nodes along with target nodes in the WSN, and derive the comparative distance relations in the relative map between possible pairs of nodes according to the RSS (Received Signal Strength) measurements using the MLE method based on the IR-UWB path loss model. The known coordinates of the anchor nodes provide the transformation matrix between the relative map and the absolute map, and the transformation matrix is used to derive the actual coordinates of the target nodes in the WSN.

NMDS-MLE is different from other MDS variants in that it only utilizes the comparative distance relationships to derive the relative map, instead of the proportional distance measurements.

Overall, the contribution of this work is the following:

- 1. A new UWB coding method, called U-BOTH (UWB based on Orthogonal Variable Spreading Factor and Time Hopping), is proposed for minimum interference communication.
- 2. An ALOHA-type channel access protocol and a message exchange protocol are used to collect distance information in WSNs.
- 3. The UWB path loss model in U-BOTH is derived and applied in the maximum likelihood estimation (MLE) method to compute the distances between neighbor nodes using the RSS information.
- 4. The NMDS-MLE (Non-metric Multidimensional Scaling and Maximum Likelihood Estimation) localization algorithm is proposed comparative distance information.

The rest of the paper is organized as follows. Section 2 describes the basic assumptions of the localization system, and the notation used in this paper. Section 3 presents a new IR-UWB coding method, called U-BOTH (UWB based on Orthogonal Variable Spreading Factor and Time Hopping), and provides the signal processing model. Section 4 specified a message exchange protocol for localization purposes using U-BOTH. According to the path loss model and the RSS information gathered by the target nodes, Section 5 and Section 6 present the ranging and localization algorithms using the MLE and NMDS methods, respectively. Section 7 evaluates the system using simulations and real testbed prototyping. Section 8 concludes the paper.

## **2 ASSUMPTIONS AND NOTATION**

We assume that a small set of sensors in a WSN work as anchors with known coordinate information, while the others are target nodes whose coordinates are to be determined. Each node in the WSN is able to exchange messages with each other using the communication protocols specified in this paper. Although our algorithms can be easily extended to carry out target tracking operations, we assume that all anchors and targets in the WSN are stationary for simplicity.

For convenience, the notation in Table 1 is used in this paper.

## **3 PHYSICAL LAYER MODEL**

#### **3.1 UWB Signal Spreading and Modulation**

In order to achieve accurate localization, a reliable physical layer communication technique is highly desirable that reduces bit error rate (BER), while mitigating the multi-users-interference (MUI) and Gaussian noise interference. Therefore, we utilize a UWB system based on time-hopping (TH) signal transmission as well as OVSF (orthogonal variable spread factor) for spreading out the symbols in the physical layer.

OVSF (OrthogonalVariable Spread Factor) was extensively used in CDMA systems to provide variable spreading codes [2]. Shorter OVSF code lengths are usually optimized for short-distance and high-data-rate transmission in less crowed environments due to its smaller spreading factor. TH (time hopping) is one of many signal modulation methods used by UWB. We describe a system, called U-BOTH (UWB modulation Based on OVSF and Time Hopping), which applies the time-hopping pulse position modulation (TH-PPM) algorithm to encode UWB pulse streams, and OVSF direct sequence to spread the user data bit stream.

Figure 1 illustrates the utilization of time hopping (TH) pulse position modulation and OVSF spreading to encode a single bit in the user data stream. First, U-BOTH sends each bit in the bit time, denoted by  $T_b$ . Then it modulates the bit 1 using a TH code, 12110021, in which each digit denotes a chip slot position within a frame time,  $T_f$ , to send a broadband radio pulse. The number



TABLE 1 Notation and Meaning

of pulses is denoted by  $N_s$ . Therefore, each bit duration is  $T_b = T_f \times N_s$ . Each chip slot lasts for  $T_c$ , sufficient to send a short UWB pulse signal.

After the initial pulse position modulation using UWB signals, the pulse sequence is again applied with OVSF code so that the phases are shifted by *π* to provide orthogonality between multiple users. The length of the OVSF code is called the spread factor *SF*, which is equal to *NS*.



FIGURE 1 U-BOTH: Interference Resistant UWB Modulation Using Time Hopping and OVSF.

In U-BOTH, the TH code is a pseudo-random sequence generated from foreknown seeds, such as node IDs. While the OVSF codes are selected from a well-defined set of orthogonal spreading codes.

To formally analyze the system in this paper, we represent the transmitted signal by the *n*th transmitter in Eq. (1):

$$
s^{n}(t) = \sum_{j=-\infty}^{+\infty} d_{j}^{n} a_{j/N_{s}j}^{n} \sqrt{E_{TX}^{n}} p_{0}(t - jT_{f} - c_{j}^{n} T_{c}),
$$
\n(1)

in which,  $d_j^n = \pm 1$  is the OVSF code with the period  $N_s$ ,  $E_{TX}^n$  is the energy of the *n*th transmitter,  $p_0(t)$  is the energy normalized pulse waveform,  $c_j^n \in$  $[0, N_c - 1]$  is the TH code with period  $N_s$  and  $a_{[j/N_s]}^n$  indicates the data stream bit. If the data bit is 1,  $a_{\lfloor j/N_s \rfloor}^n = +1$ . Otherwise,  $a_{\lfloor j/N_s \rfloor}^n = -1$ .

At the receiver side, the received signal consists three source of information:

$$
r(t) = r_u(t) + r_{mui}(t) + n(t),
$$

in which,  $r_u(t)$  is the desired user signal,  $r_{mu}(t)$  is co-channel interference from multiple users, and  $n(t)$  is the additive white Gaussian noise (AWGN).

Denote the pulse energy of the *n*-th transmitter as  $E_{RX}^n$ . Without loss of generality, we assume that the first user's transmission is the desired signal at the receiver for simplicity, then Eq. (2) provides the desired signal function at the receiver:

$$
r_u(t) = \sum_{j=-\infty}^{+\infty} d_j^1 a_{j/N_s}^1 \sqrt{E_{RX}^1} p_0(t - jT_f - c_j^1 T_c).
$$
 (2)

We define the correlation template of the receiver:

$$
m(t) = \sum_{j=iN_s}^{(i+1)N_s-1} d_j^1 p_0(t - jT_f - c_j^1 T_c); i \in (-\infty, +\infty).
$$
 (3)

# **3.2 Single User System Analysis**

As the first step, we assume that the channel isAWGN multipath-free channel, and that the transmitter and the receiver are synchronized. In a single user signal processing system, the input of the receiver has two parts:  $r<sub>u</sub>(t)$  and  $n(t)$ , and the output of the receiver in time interval  $[0, T_b]$  is represented by:

$$
Z = Z_u + Z_n = \int_0^{T_b} (r_u(t) + n(t))m(t)dt.
$$
 (4)

In Eq. (4), the useful output signal is:

$$
Z_u = \sum_{j=0}^{N_s-1} \int_{jT_f+c_j^1T_c}^{jT_f+c_j^1T_c} d_j^1 d_j^1 a_{j/N_s}^1 \sqrt{E_{RX}^1} \omega(t) dt,
$$

where  $\omega(t) = p_0(t - jT_f - c_j^1 T_c)p_0(t - jT_f - c_j^1 T_c)$ .

Because  $d_j^1 d_j^1 = 1$ ,  $p_0(t)$  is the energy normalized pulse waveform, we have

$$
Z_u = \sum_{j=0}^{N_s - 1} \int_0^{T_c} a_{\lfloor j/N_s \rfloor}^1 \sqrt{E_{RX}^1} p_0(t) p_0(t) dt
$$
  
=  $N_s a_{\lfloor j/N_s \rfloor}^1 \sqrt{E_{RX}^1} \int_0^{T_c} p_0(t) p_0(t) dt$   
=  $a_{\lfloor j/N_s \rfloor}^1 N_s \sqrt{E_{RX}^1}$ 

In Eq. (4), the output noise signal is:

$$
Z_n = \sum_{j=0}^{N_s-1} \int_0^{T_c} d_j^1 p_0(t) n(t) dt = \sum_{j=0}^{N_s-1} d_j^1 n_j,
$$

where  $n_j$  is Gaussian random variable with mean 0 and variance  $N_0/2$ . Because  $d_j^1$  is not a random variable, the variance of  $Z_n$  is:

$$
D(Z_n) = D(\sum_{j=0}^{N_s-1} d_j^1 n_j) = N_s \frac{N_0}{2},
$$
  

$$
Z_n \sim N(0, N_0 N_s/2).
$$

Suppose that the statistical probabilities of data bit  $b = 0$  and  $b = 1$  are equal, we obtain the BER (bit error rate) of the single user system in AWGN channel as follows:

$$
Pr_b = \frac{1}{2}P(Z > 0|b = 0) + \frac{1}{2}P(Z < 0|b = 1) = P(Z > 0|b = 0).
$$

Because  $a_{[j/N_s]}^1 = -1$  if  $b = 0$ , then the useful output is  $Z_u = a_{[j/N_s]}^1 N_s \sqrt{E_{RX}^1} =$  $-N_s\sqrt{E_{RX}^1}$ . Using Eq. (4), the BER become:

$$
Pr_b = P(Z > 0 | b = 0) = P(-N_s \sqrt{E_{RX}^1} + Z_n > 0)
$$
  
=  $P(Z_n > N_s \sqrt{E_{RX}^1})$ 

It can be rewritten by complementary error function erfc  $(x)$  as follow:

$$
Pr_b = \frac{1}{2} \operatorname{erfc}\left(\sqrt{\frac{N_s E_{RX}^1}{N_0}}\right).
$$

Where  $erfc(x) = \frac{2}{\sqrt{2}}$  $\frac{2}{\prod} \int_{x}^{\infty} \exp(-t^2) dt.$ 

Because U-BOTH is a rate variable system using OVSF, we analyze the relation between BER and the bit rate. Suppose the system's OVSF code is a code tree of 6 layers [6], and the spreading factor is 2, 4, 8, 16, 32, 64, respectively. Further suppose the basic rate of our system is  $R_0$ , then the corresponding bit rate of U-BOTH is  $R_b = iR_0$  (*i*= 32, 16, 8, 4, 2, 1, respectively).

Denote the bit rate as  $R_b$ , where  $R_b = iR_0$ ,  $i = 1, 2, \dots, 32$ , we can get the relation between BER and the bit rate:

$$
Pr_b = \frac{1}{2} \operatorname{erfc} \left( \sqrt{\frac{SF \cdot E_{RX}^1}{N_0}} \right)
$$
  
=  $\frac{1}{2} \operatorname{erfc} \left( \sqrt{\frac{64R_0 \cdot E_{RX}^1}{R_b N_0}} \right)$  (5)

Eq. (5) shows that the BER decrease when the spreading factor *SF* increases or when the bit rate decreases. Therefore, we can adjust *SF* to adapt different environments with various noise levels while maintaining the same bandwidth of the signal. This is the main reason we adjust OVSF codes in our system.

#### **3.3 Multi-User Interference Analysis**

In multi-user communication system, the received signal includes multi-user interference  $Z_{mui}$  and noises. The  $Z_u + Z_n$  part is the same as Eq. (4), but the



FIGURE 2 The Interference to User 1 by The *n*-th User.

multi-user interference  $Z_{mui}$  is additional. Because the phase and delay  $\tau$  of interfering pulses is random as shown in Figure 2, we have to compute the interference's variance.

Suppose that  $\tau^n$  is uniformly distributed over [0,  $T_f$ ), then the interference variance of the desired signal, *i.e.* the signal from the 1st user, caused by transmitter  $n$  is [13]:

$$
\sigma_{bit}^2 = \frac{N_s}{T_f} \int_0^{T_f} \left( \sqrt{E_{RX}^n} \int_0^{T_c} d_j^1 d_i^n p_0(t - \tau^n) p_0(t) dt \right)^2 d\tau^n.
$$

Therefore, the total interference variance  $\sigma_{mui}^2$  from all other transmitters is:

$$
\sum_{n=2}^{N_u} \left( \frac{N_s E_{RX}^n}{T_f} \int_0^{T_f} \left( \int_0^{T_c} d_j^1 d_i^n p_0(t-\tau^n) p_0(t) dt \right)^2 d\tau^n \right).
$$

Because the delay  $\tau$  for all transmitters has the same distribution, we get the following formula:

$$
\sigma_{mui}^2 = \frac{N_s}{T_f} \sum_{n=2}^{N_u} E_{RX}^n \left( \int_0^{T_f} \left( \int_0^{T_c} d_j^1 d_i^n p_0(t - \tau^n) p_0(t) dt \right)^2 d\tau^n \right)
$$
  
=  $\sigma_M^2 \frac{N_s}{T_f} \sum_{n=2}^{N_u} E_{RX}^n$ 

in which,

$$
\sigma_M^2 = \int_0^{T_f} \left( \int_0^{T_c} d_j^1 d_i^n p_0(t - \tau^n) p_0(t) dt \right)^2 d\tau
$$
  
= 
$$
\int_0^{T_f} R^2(\tau) d\tau.
$$

According to [13], and noticing that  $R_b = \frac{1}{N_s N_f}$  and  $N_s = SF = \frac{64R_0}{R_b}$ , Eq. (6) gives the BER in multi-user interference environments.

$$
Pr_b = \frac{1}{2} \operatorname{erfc} \left( \sqrt{\frac{1}{2} \left( \left( \frac{2N_s E_{RX}^1}{N_0} \right)^{-1} + \left( \frac{N_s E_{RX}^1}{\sigma_M^2 \frac{1}{T_f} \sum_{n=2}^{N_u} E_{RX}^n} \right)^{-1} \right)} \right) = \frac{1}{2} \operatorname{erfc} \left( \sqrt{\frac{1}{2} \left( \left( \frac{128R_0 E_{RX}^1}{R_b N_0} \right)^{-1} + \left( \frac{E_{RX}^1}{\sigma_M^2 R_b \sum_{n=2}^{N_u} E_{RX}^n} \right)^{-1} \right)^{-1}} \right).
$$
 (6)

## **4 NETWORK PROTOCOL OPERATIONS**

U-BOTH based communication requires each node to encode signals using specific OVSF codes. In distributed wireless networks, code assignments are categorized into transmitter-oriented, receiver-oriented or a per-link-oriented code assignment schemes (also known as TOCA, ROCA and POCA, respectively) [10, 12]. Regarding the ways of assigning the OVSF-TH codes and encoding the MAC data frames for transmissions, we propose two different ways to implement multiple access protocols using U-BOTH.

**ROCA-Based Protocol Operations** The first approach is based on the receiver-oriented code assignment (ROCA), in which case the data packet transmissions are encoded using the unique OVSF-TH code assigned to the receiver. Beside ROCA, there is a common OVSF-TH code for bootstrapping and coordination purposes.

In ROCA scheme, in order to carry out localization algorithms, each node sends location request messages using the common OVSF-TH code to the neighbor nodes. The request message includes the request command, and the receiver's OVSF-TH code. Upon receiving the request message, neighbor nodes sends back a response message using the receiver's OVSF-TH code using a random backoff mechanism. Handshake messages provide both RSS information between pairs and the estimated coordinates in the current rounds.

**TOCA-Based Protocol Operations** The second approach is based on transmitter-oriented code assignment (TOCA), in which case each packet transmission is encoded using two OVSF-TH codes — one is a common OVSF-TH code to encode the common physical layer frame header, and the other transmitter-specific code is to encode the physical layer frame payload. The frame head includes the transmitter-oriented OVSF-TH code for encoding the frame payload.

Because the physical layer headers are sent on a common OVSF-TH code, the physical layer header transmissions resemble those of ALOHA networks with regard to packet collision. Because the headers are usually short, the collision probability is low.

On the other hand, because the data frame payload is transmitted on unique OVSF-TH codes, the interference between the payload and other frame headers and payloads is dramatically reduced.

In both ROCA- and TOCA-based systems, packets from the neighbor nodes can be lost. However, this does not affect the overall performance of our localization algorithms because they tolerate such losses.

After getting the respective RSS information between each pair of nodes, the network calculates its coordinate in two steps — ranging and localization.

## **5 RANGING ALGORITHM**

Ranging is to estimate the approximate distance between adjacent nodes. We use the MLE (maximum likelihood estimation) method for such calculations. First of all, we need to establish the path loss model of the UWB channel in order to inversely derive the distance information from received signal qualities.

## **5.1 The Path Loss Model**

It is well-known that the path loss model can be expressed by the log-distance path loss law in many indoor or outdoor environments, as shown by Eq. (7).

$$
PL(d) = \left( PL_0 + 10\gamma \log_{10}(\frac{d}{d_0}) \right) + S; \ d \ge d_0,
$$
 (7)

in which

- $\bullet$  *d*<sub>0</sub> is the reference distance (e.g. 1 meter in UWB medium),
- $PL_0$  means the path loss in dB at  $d_0$ ,
- *d* is the distance between the transmitter  $(Tx)$  and receiver  $(Rx)$ ,
- *γ* refers to the path loss exponent which depends on channel and environment,
- *S* is the log-normal shadow fading in dB. Usually, *S* is a Gaussiandistributed random variable with zero mean and standard deviation *σS*.

Eq. (7) could construct a statistical path loss model for UWB propagation in different environments. The path loss *PL*(*d*) can be expressed as a Gaussiandistributed random variable with:

$$
S \sim N(0, \sigma_S^2),
$$

$$
PL(d) \sim N(PL_0 + 10\gamma \log_{10} d, \sigma_S^2).
$$

The probability density function (*pdf*) of path loss *PL*(*d*) is:

$$
p(PL) = \frac{e^{-\frac{[PL - (PL_0 + 10\gamma \log_{10} d)]^2}{2\sigma_S^2}}}{\sqrt{2\pi \sigma_S^2}}.
$$
\n(8)

IEEE 802.15.4a Task Group provided Channel Model 1-9 by taking limited real measurements to determine the values of  $\gamma$ ,  $\sigma_s$  and other variables in different situations. When deploying real UWB networks, people could approximately choose the corresponding channel model with the parameters specified in IEEE 802.15.4a.

#### **5.2 Ranging Algorithm based on Maximum Likelihood Estimation**

The distance between the transmitter  $Tx$  and the receiver  $Rx$  in Eq. (7) can be calculated by the general ranging method between two nodes using the RSS information:

$$
\hat{d} = 10^{\frac{PL(d) - PL_0 - S}{10\gamma}}.
$$

Receiver computes the distance between the transmitter *Tx* and the receiver *Rx* using random values *S*. However, in above single random ranging, the random variables *S* selected by the sensor nodes are not exactly those in the real time-variant channel. In order to avoid the ranging errors caused by the large deviation between the estimated *S* values and the real *S* values in each round of ranging estimation, we propose an iterative ranging based on MLE (maximum likelihood estimation) in UWB wireless sensor networks.

Suppose *PLi* is the *i*th observation value, we get the joint conditional *pdf*  $p(PL|d)$  using Eq. (9).

$$
p(PL|d) = \prod_{i=1}^{N} \frac{e^{-\frac{[PL_i - (PL_0 + 10\gamma \log_{10} d)]^2}{2\sigma_S^2}}}{\sqrt{2\pi \sigma_S^2}}.
$$
(9)

The necessary condition to compute the MLE of *d* is:

$$
\frac{\partial \ln p(PL|d)}{\partial d} = \frac{10N\gamma}{\sigma_S^2 d \ln 10} \left( \frac{1}{N} \sum_{i=1}^N PL_i - PL_0 - 10\gamma \log_{10} d \right)
$$
  
= 0. (10)

We solve Eq.  $(10)$  and have:

$$
\widehat{\log_{10} d} = \frac{1}{10N\gamma} \sum_{i=1}^{N} PL_i - \frac{PL_0}{10\gamma}
$$

Therefore, the MLE based RSS UWB ranging is:

$$
\hat{d} = 10^{\frac{1}{10N\gamma} \sum_{i=1}^{N} PL_i - \frac{PL_0}{10\gamma}}.
$$
\n(11)

## **6 LOCALIZATION ALGORITHM**

#### **6.1 Multi-Dimensional Scaling (MDS)**

MDS (Multidimensional Scaling) is a statistical technique for exploratory data analysis or information visualization. MDS collects the proximity data between each pair of spatial objects as reference. Then it visualizes objects as points in a low dimensional Euclidean space and represents these proximity data as distances between points. In order to derive accurate results, MDS has to find some solutions that relate distance information to proximity information as closely as possible.

Suppose that *n* denotes the number of different objects, and the proximity for objects *i* and *j* is denoted by  $p_{ii}$ . Thus, we derive a proximity matrix  ${\bf P}_{n \times n} = p_{ii}$ . The coordinates of mapping points are represented by a matrix  $\mathbf{X}_{n \times m}$ , where *m* is the dimensions of the solution, e.g. 2D or 3D.

Now, let  $d_{ii}(\mathbf{X})$  be the Euclidean distance between points *i* and *j* with coordinates in  $\mathbf{X}_{n \times m}$ , respectively. The objective of MDS is to find a matrix **X** so that  $d_{ii}(\mathbf{X})$  proportionally matches  $p_{ii}$  as closely as possible, which is presented by  $f(p_{ii}) \sim d_{ii}(\mathbf{X})$ . The closeness is measured by metric *STRESS* as follows:

$$
STRESS = \sum [f(p_{ij}) - d_{ij}(\mathbf{X})]^2.
$$

MDS algorithms are categorized into several types, depending on whether the similarity data is quantitative or qualitative, and are called metric MDS and non-metric MDS, respectively.

Classical metric MDS formulates the relationship between proximity data of objects and distances in the Euclidean space by transformation functions. In order to find a perfect fitness between proximity data and Euclidean distance, the transformation formula  $d_{ii}(\mathbf{X}) = f(p_{ii})$  is pursued, such as a linear model:  $d_{ii}(\mathbf{X}) = a + bp_{ii}$ . Because  $d_{ii}(\mathbf{X})$  represents the Euclidean distance between points*i* and *j* in coordinate matrix **X**, MDS rests on the fact that the coordinate matrix  $X$  can be derived by double centering and eigenvalue decomposition from the proximity matrix **P** with the least error.

The relationship between the proximity of objects and the Euclidean distances of points in Non-metric MDS is not as strict as metric MDS. Non-metric MDS only requires a monotonic relationship between them.

When Non-metric MDS takes proximity data of different objects to construct corresponding spatial coordinates, it only requires that the rank order of the proximity  $p_{ij}$  have to keep the same ordinal level as the distances  $d_{ij}$ . That is,

$$
\forall i,j,k,l: p_{ij} < p_{kl} \Rightarrow d_{ij}(X) < d_{kl}(X).
$$

Compared with metric MDS, the monotonic assumption that the data is measured at the ordinal level in Non-metric MDS makes it more flexible and applicable for localization in wireless sensor networks.

## **6.2 The NMDS-MLE Localization Algorithm**

NMDS-MLE localization algorithm combines the ranging and localization processes. Ranging is based on the iterative RSS information collected by above U-BOTH UWB system and refined by the MLE method. Localization is based on the NMDS algorithm. As a whole, NMDS-MLE localization consist of 5 steps:

- 1 Gather RSS information between neighbors by U-BOTH system in the network, and form a sparse matrix **R**, which is derived from the estimated distances denoted by  $r_{ii}$ .  $r_{ii}$  is estimated by iterative ranging based on RSS information and MLE method. For the nodes that is out of the communication range, *rij* is zero.
- 2 Construct the proximity data matrix **P** based on sparse matrix **R**. The estimated distance  $p_{ij}$  between every pair of nodes in the network is computed by the shortest path algorithm, such as the Dijkstra's or Floyd's algorithm.
- 3 Construct the coordinate system to plot the objects in the Euclidean space and obtained the distance matrix *D* composed by the Euclidean distance *dij*.
- 4 Compare the ordinal level between aforementioned two types of distance information: estimated distance  $p_{ij}$  and Euclidean distance  $d_{ij}$ , and refine the relative coordinate **X** of nodes in Non-metric MDS.
- 5 Transform relative coordinate into global absolute coordinate by the anchor nodes in the network.

In Step 3 and Step 4, localization is executed by NMDS-MLE as Algorithm 1. In Algorithm 1, a monotonic transformation between proximity data and Euclidean distance is calculated in line 6 to 14, which yields an intermediate distance value  $\hat{d}_{ij}$ . By performing a monotone regression with the current distances  $d_{ij}$  as targets and proximity  $p_{ij}$  as inputs, NMDS-MLE generates  $d_{ij}$  to reflect the ordinal level of  $p_{ij}$  in each iteration, where  $\hat{d}_{ij}$  should be subjected to:

**Input**: node set *N*, initial coordinate matrix  $X^{(0)}$ , proximity data matrix *P*, threshold  $\varepsilon$ , iteration number *k* ←− 0 **Output**: relative coordinate  $X^{(n)}$ **1 for** *each*  $i, j \in N$  **do** 2  $\int d^k_{ij}$  ←  $\sqrt{(x^k_i - x^k_j)^2 + (y^k_i - y^k_j)^2}$ 3  $\cos$  construct the Euclidean distance matrix  $D^{(k)}$ **4 end 5 while**  $STRESS \geq \varepsilon$  **do 6 for** *each i*, *j*, *u*, *v*  $\in$  *N* **do**<br>**7 if**  $p_{ii} < p_{uv}$  **and**  $d_{ii}^k$ *7*  $\vert$  **if**  $p_{ij} < p_{uv}$  and  $d_{ij}^k > d_{uv}^k$  then **8**  $\begin{vmatrix} \hat{d}_i^k \leftarrow (d_{ij}^k + d_{uv}^k)/2 \end{vmatrix}$ **9**  $\begin{vmatrix} \hat{d}_{uv}^k \longleftarrow (d_{ij}^k + d_{uv}^k)/2 \end{vmatrix}$ **10**  $\qquad$  **else if**  $p_{ij} < p_{uv}$  and  $d_{ij}^k \leq d_{uv}^k$  then  $\hat{d}^k_{ij} \longleftarrow d^k_{ij}$  $\hat{d}^k_{uv} \leftarrow d^k_{uv}$ **13 end 14 end 15**  $k \leftarrow k+1$ <br>**16** undate the c update the coordinate matrix  $X^{(k)}$ 17  $\Box$  update the distance matrix  $D^{(k)}$ **18 end**

$$
\forall i, j, k, l : p_{ij} < p_{kl} \Rightarrow \hat{d}_{ij}(X) < \hat{d}_{kl}(X).
$$

Because of above relation between  $p_{ij}$  and  $\hat{d}_{ij}$ , NMDS-MLE takes following *STRESS* applied in line 5 to evaluate the accuracy of the fitting:

$$
STRESS = \sqrt{\sum_{ij, i \neq j} (\hat{d}_{ij} - d_{ij})^2 / \sum_{ij, i \neq j} d_{ij}^2}.
$$
 (12)

A small *STRESS* indicates a good fit, whereas a high value indicates a bad fit. Kruskal [11] provide some guide lines of stress value with respect to the goodness of fit of the solution, shown in Table 2.

Note in line 16 and line 17 in Algorithm 1, NMDS-MLE updates spatial coordinate matrix  $X^{k-1}$  to  $X^k$  according to  $d_{ij}^{k-1}$  and  $\hat{d}_{ij}^{k-1}$ , and then obtains a new Euclidean distance  $d_{ij}(\mathbf{X})^k$ . The spatial coordinate  $(x_i^k, y_i^k)$  is updated as follows:

$$
x_i^k = x_i^{k-1} + \frac{\alpha}{n-1} \sum_{j \in M, j \neq i} \left( 1 - \frac{\hat{d}_{ij}^{k-1}}{d_{ij}^{k-1}} \right) (x_j^{k-1} - x_i^{k-1}),
$$
  
\n
$$
y_i^k = y_i^{k-1} + \frac{\alpha}{n-1} \sum_{j \in M, j \neq i} \left( 1 - \frac{\hat{d}_{ij}^{k-1}}{d_{ij}^{k-1}} \right) (y_j^{k-1} - y_i^{k-1}).
$$
\n(13)



#### TABLE 2 Stress and goodness of fit

Where *n* is the number of target nodes,  $\alpha$  is the iteration step length, which is set to be 0.2 in the paper.

In Step 4, the estimated location matrix **X** represents the relative coordinates of nodes, which have a different orientation and scaling than the original coordinates. And in Step 5, the transformation from relative coordinate **X** into absolute coordinates usually includes shift, rotation, scaling, and reflection of coordinates, which are implemented by some transformation to minimize the errors between the absolute coordinates of anchor nodes and their relative locations in the NMDS map. Suppose there are *m* anchor nodes whose relative locations are  $X_R = (X_{R_1}, X_{R_2}, \cdots, X_{R_m})$ , and real locations are  $X_T = (X_{T_1}, X_{T_2}, \cdots, X_{T_m})$ . First, we derive optimal transformation function **Q**, then transfer all the relative coordinates of nodes to the absolute coordinates by the optimal transformation function **Q**.

## **7 EVALUATIONS**

Our algorithms were implemented in Matlab 7.0 to verify our localization algorithms based on U-BOTH system for WSNs. The datum for evaluations were collected and validated through two sets of experiments, one based on simulations and the other on real experiments.

In the simulations, we set the following scenarios for evaluations:

- 1. With regard to the BER (bit error rate), we evaluated U-BOTH system performance in single and multi-user scenarios.
- 2. Using NMDS-MLE localization algorithm, we evaluated our localization model both in random network and grid network.

In the real testbed experiments, we deployed wireless sensors that run the IEEE 802.11 standard. The mathematical tools and the results of these experiments similarly applied to other wireless systems. In our testbed, we deployed five WiFi nodes, with the same transmission power (30dBm) on our office floor as represented by the solid dots with ID labels in Figure 3. The transmission radius of the WiFi nodes is approximately 30 meters.



FIGURE 3 Testbed Experiment Scenario.



## FIGURE 4

Parameters on Log-distance Path Loss Model.

In both simulated and testbed experiments, we adopted the log-distance UWB path loss model in NLOS (Non-Line-of-Sight) environment for our estimators, which is parameterized in Figure 4.

We evaluated the performance of localization algorithms with mean estimation error, which was widely used in previous research works:

$$
error = \frac{\sum_{i=m+1}^{n} \|X_{est}^{i} - X_{real}^{i}\|^2}{(n-m) \times R} \times 100\%
$$
\n(14)

where *n* and *m* are the total number of sensors and the number of anchor nodes in the WNS, respectively, *R* represents communication range.

## **7.1 Simulation Results**

## *U-BOTH System Performance*

We assumed the channel is AWGN multipath-free single user channel, the transmitter and the receiver are synchronized perfectly. Then we randomly generated 2000 bits, every bit uses 4 pulses to repeat coding  $(N_s = 4)$ .

Figure 5 illustrates the BER of the received signal using U-BOTH system, in contrast to DS-UWB that only uses direct sequence spreading, and TH-UWB that uses time-hopping pulse position modulation alone for UWB transmissions. We can see that the BER of U-BOTH and the DS-UWB system which use the  $\pi$ -phase shift keying modulation are lower than TH-UWB. This is because the distance of two signals in binary phase shift keying (BPSK) modulation is  $2\sqrt{E_{pulse}}$ , but  $\sqrt{2E_{pulse}}$  in TH-UWB [17].

Secondly, we let  $E_b/N_0 = 0$  dB,  $N_s = 4$  and generated 2000 bits randomly. Figure 6 shows the relative performance of U-BOTH, TH-UWB and DS-UWB systems in multiple access scenarios. In this case, the received signal includes noise and co-channel interference. In Figure 6, although both the BER and the variance of error bits increase as the number of users increases, the performance of our U-BOTH system is still better than DS-UWB and TH-UWB, proving that the UWB coding based OVSF-TH effectively handle the burst errors.

## *Evaluation of the Localization Algorithms*

**Random Deployment** 100 nodes were deployed randomly in a 100*m* × 100*m* square area as shown in Figure 7(a), in which points represent nodes and edges represent the connections between neighbor nodes. The communication



FIGURE 5 Bit Error Rate in A Single User System with Additive White Gaussian Noise (AWGN).



FIGURE 6 Bit Error Rate and The Variance of The Number of Error Bits of 2000 Generated Bits.

range is 12 *m* and the average connectivity is 4.6. The *STRESS* value was set between .00 and .025 guide line.

Figure 7(b) reflects the relative coordinate of every node generated by NMDS-MLE. It shows that the relative coordinates have a different orientation and scaling than the original network in Figure  $7(a)$ . This is because that relative coordinate is derived only based on the distance relationship between every pair of nodes in the network. Figure 7(c) derives the absolute coordinates of all the nodes. Their relative coordinates in Figure 7(b) are transformed based on the location information provided by 4 random anchor nodes denoted by  $\times$ . The dots represent the real locations of the nodes, and the lines with arrows indicate the errors of the estimated locations from the real locations, the average localization error is about 5.3850%. The MDS-MAP algorithm is also applied in the case and the average localization error is about 18.4747%.

**Grid Deployment** 100 nodes were deployed in a  $45m \times 45m$  square area with grid deployment in Figure 8(a). The communication range is 12 *m* and the average connectivity is 16.8. The *STRESS* value was set between .00 and .025 guide line.With the same symbol meaning in the figures, Figure 8(b) represents the relative coordinate map using NMDS-MLE algorithm and Figure 8(c) depicts the absolute coordinate map by transformation based on 4 random anchor nodes. The average localization error in the grid case is about 1.2876%. For MDS-MAP algorithm, it is about 6.0543%.

#### *Performance analysis*

The localization performance of NMDS-MLE in different scenarios under different degrees of connectivity is analyzed by Figure 9, compared with the MDS-MAP by the same experimental settings. From the figure, we can see that localization error of NMDS-MLE algorithm is much lower and more



(c) Absolute Coordinates

FIGURE 7 Random Deployment.



(a) Original Coordinates



#### (b) Relative Coordinates



(c) Absolute Coordinates

FIGURE 8 Grid Deployment.



FIGURE 9 Relation between The Connectivity and The Localization Error.

stable than MDS-MAP in different scenarios. Furthermore, when NMDS-RSS and MDS-MAP are applied in grid deployment with varies of connectivity. It shows that NMDS-RSS obtain higher localization accuracy in the grid layout than in the random layout for the same connectivity level.

Figure 10 presents the relation between localization error and the number of iteration *N* in NMDS-MLE algorithm. Because the accuracy of ranging is improved by MLE method based on the RSS information provided by our U-BOTH system, it is obvious that the localization error decreases dramatically when the number of iterations in ranging increases in both random and grid deployment.



FIGURE 10 Relation between The Number of Iteration *N* and The Localization Error.

## *Testbed Experimental Results*

In the testbed experiments, we evaluated the real effects of NMDS-MLE localization algorithm. Five nodes were deployed on our office floor under the NLOS environment as shown in in Figure 3. The communication range is 30 *m* and the connectivity is 5. The *STRESS* value was set below .00 guide line. We randomly selected node 1 and node 3 as the anchor nodes.

With the same symbol meaning in the figures, Figure 11(a) represents the absolute coordinate map by collecting 20 RSS values for iterations in NMDS-MLE algorithm. The average localization error is about 2.1256%. Meanwhile, Figure 11(b) represents the localization results by collecting 100 RSS values, and the corresponding localization error is about 0.7522%.



FIGURE 11 Testbed Experiment Estimation Results of NMDS-MLE.

From the real experiments, it shows that given proper *STRESS* value and iteration number, NMDS-MLE performs fairly accurately for location estimation under complex signal propagation environments.

## **8 CONCLUSION**

In order to provide a localization algorithm using the NMDS-MLE methods, we have proposed the communication protocols based on a new UWB coding method, called U-BOTH (UWB based on Orthogonal Variable Spreading Factor and Time Hopping), and an ALOHA-type channel access method and a message exchange protocol to collect distance information in WSNs. Then we specified the NMDS-MLE algorithms using the UWB path loss model for ranging and localization purposes. The performance of NMDS-MLE algorithms in the U-BOTH based communication system were analyzed using communication theories and simulations. Simulation and real testbed experiment results show that U-BOTH transmission technique can effectively reduce the bit error rate under the path loss model, and the corresponding ranging and localization algorithms can achieve comparable or better results than previous localization methods.

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